Learning Dynamics and Equilibrium Selection

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Networked interaction: Societal, engineered, & hybrid







• How could agents converge to NE?

Arrow: "The attainment of equilibrium requires a disequilibrium process."

- Monographs:
 - Weibull, Evolutionary Game Theory, 1997.
 - Young, Individual Strategy and Social Structure, 1998.
 - Fudenberg & Levine, The Theory of Learning in Games, 1998.
 - Samuelson, Evolutionary Games and Equilibrium Selection, 1998.
 - Young, Strategic Learning and Its Limits, 2004.
 - Sandholm, Population Dynamics and Evolutionary Games, 2010.
- Surveys:
 - Hart, "Adaptive heuristics", *Econometrica*, 2005.
 - Fudenberg & Levine, "Learning and equilibrium", Annual Review of Economics, 2009.

- If agents self-organize to Nash equilibrium...
 - Price of Anarchy:



	А	В			S	Н
А	4,4	0,0		S	3,3	0,1
В	0,0	3,3		H	1,0	1,1
Typewriter Game					Stag	Hunt

- How to distinguish equilbria?
- Payoff based distinctions: Payoff dominance vs Risk dominance
- Evolutionary (i.e., *dynamic*) distinction
 - Young (1993) "The evolution of convention"
 - Kandori/Mailath/Rob (1993) "Learning, mutation, and long-run equilibria in games"
 - many more...
- Adaptive play:
 - "Two" players sparsely sample from finite history
 - Players either:
 - * Play best response to selection
 - * Experiment with small probability
 - Young (1993): Risk dominance is "stochastically stable"

- Dynamics & equilibrium selection theme continued...
 - Constrained log linear learning
 - Self assembly
 - Dynamic reinforcement dynamics
- "Prescriptive" issues & opportunities
 - What are implications of additional constraints?
 - How to exploit additional degrees of freedom?

- Setup:
 - Players: $\{1,...,p\}$
 - Actions: $a_i \in \mathcal{A}_i$
 - Action profiles:

$$(a_1, a_2, ..., a_p) \in \mathcal{A} = \mathcal{A}_1 \times \mathcal{A}_2 \times ... \times \mathcal{A}_p$$

- Payoffs: $u_i : (a_1, a_2, ..., a_p) = (a_i, a_{-i}) \mapsto \mathbf{R}$
- Global objective: $G : \mathcal{A} \rightarrow \mathbf{R}$
- Action profile $a^* \in A$ is a *Nash equilibrium* (NE) if for all players:

$$u_i(a_1^*, a_2^*, ..., a_p^*) = u_i(a_i^*, a_{-i}^*) \ge u_i(a_i', a_{-i}^*)$$

- Learning dynamics:
 - $-t = 0, 1, 2, \dots$
 - $-\mathbf{Pr}[a_i(t)] = p_i(t), \quad p_i(t) \in \Delta(\mathcal{A}_i)$
 - $p_i(t) = \mathcal{F}_i(available info at time t)$

• Potential games: For some $\phi : \mathcal{A} \to \mathbb{R}$

$$\phi(a_i, a_{-i}) - \phi(a'_i, a_{-i}) > 0$$
$$\Leftrightarrow$$
$$u_i(a_i, a_{-i}) - u_i(a'_i, a_{-i}) > 0$$

i.e., potential function increases iff unilateral improvement.

• Features:

- Typical of "coordination games"
- Desirable convergence properties under various algorithms
- Need not imply "cooperation" or $\phi = G$
- Prescriptive opportunity: Potential game by design

Illustrations

- Distributed routing
 - Payoff = negative congestion. $c_r(\sigma_r)$
 - Potential function:

$$\phi = \sum_{r} \sum_{n=1}^{\sigma_r} c_r(n)$$

- Overall congestion:

$$G = \sum_{r} \sigma_r c_r(\sigma_r)$$

- Note: $\phi \neq G$
- Multiagent sudoku:

 $u_i(a) =$ #reps in row + #reps in column + #reps in sector

$$\phi(a) = \sum_i u_i(a)$$



• Preliminary: Gibbs distribution

 $\mathbf{Pr}\left[v_{i}\right] \propto e^{v_{i}/T}$

As $T \downarrow 0$ assigns all probability to $\arg \max \{v_1, v_2, ..., v_n\}$

• At stage t

- Player *i* is selected at random
- Chosen player selects

$$\mathbf{Pr}\left[a_{i}(t)=j\right] \propto e^{u_{i}(j,a_{-i}(t-1))/T}$$

- Interpretation: Noisy best reply to previous joint actions
- Fact: SAP results in a Markov chain over joint action space A with a unique stationary distribution, μ .
- Blume (1993): In (cardinal) potential games, steady-state distribution satisfies

$$\mathbf{Pr}\left[a
ight] \propto e^{\phi\left(a
ight)/T}$$

- Implication: As $T \downarrow 0$, concentrated at potential maximizer
- i.e., Potential maximizer is "stochastically stable"



• Impose constrained evolution:

$$a_i(t) \in C(a_i(t-1))$$

- Limited mobility
- Obstacles
- Mimicking log linear learning *alters* stochastically stable action profiles!
- Example: Identical interest game

$$C_2(L) = \{L, M\}$$
 $C_2(M) = \{L, M, R\}$ $C_2(R) = \{M, R\}$

- Potential maximizer: (D, L)
- Stochastically stable state: (U, R)
- Intuition:
 - $* (U, R) \rightarrow (D, L)$ "costs" 18 (used to cost 9)
 - $* (D, L) \rightarrow (U, R)$ "costs" 10 (used to cost 10)

- At stage *t*:
 - Player i is selected at random
 - Chosen player compares $a_i(t-1)$ with randomly selected $a'_i \in C(a_i(t-1))$

 $\mathbf{Pr}\left[a_{i}(t)\right] \propto e^{u_{i}(a_{i}(t-1),a_{-i}(t-1))/T} \quad \mathbf{VS} \quad e^{u_{i}(a_{i}',a_{-i}(t-1))/T}$

- Marden & JSS, 2008: Under binary log linear learning, only potential function maximizers are stochastically stable.
- *No longer* characterize stationary distribution.
- Recall example:

	L	Μ	R
U	0	0	9
D	10	-10	-10

• Binary version: $(U, M) \rightarrow (U, L)$ now has zero resistance.

- What if evaluation of $U_i(a', a_{-i}(t-1))$ no longer possible?
- New setup: Players can only measure $a_i(t)$ and $U_i(a(t))$
- Introduce *baseline action* $a_i^b(t)$ and *baseline utility* $u_i^b(t)$
- Action selection:

 $a_i(t) = a_i^b(t)$ with probability $(1 - \epsilon)$

 $a_i(t)$ is chosen randomly over \mathcal{A}_i with probability ϵ

• Baseline action & utility update:

New baseline
with probability $\sim e^{U_i(a(t))/T}$ Keep baseline
with probability $\sim e^{u_i^b(t)/T}$ \Downarrow \Downarrow \downarrow \Downarrow $a_i^b(t+1) = a_i(t)$ $a_i^b(t+1) = a_i^b(t)$ $u_i^b(t+1) = u_i(a(t))$ $u_i^b(t+1) = u_i^b(t)$

Marden & JSS, 2008: Under payoff based linear learning, only potential function maximizers are stochastically stable.

- Definition:
 - Let P^{ϵ} denote the transition probability matrix of an irreducible & aperiodic Markov chain.
 - Let μ^{ϵ} be the (unique) stationary distribution for P^{ϵ}
 - A state, x, is stochastically stable if

$$\liminf_{\epsilon \to 0} \mu^{\epsilon}(x) > 0$$

- Young (1993): To determine stochastic stability
 - View learning dynamics as ϵ perturbation of reference ($\epsilon = 0$) Markov chain
 - Divide reference Markov chain into recurrence classes
 - Define *resistance* to transition between recurrence classes:

$$0 < \lim_{\epsilon \downarrow 0} \frac{P_{ij}^{\epsilon}}{\epsilon^{r(i \rightarrow j)}} < \infty$$

- Form *stochastic potential* for each recurrence class
- Minimal stochastic potential implies stochastic stability



Combinatoric utilization vs pragmatic utilization

Illustration: Sensor allocation

• Objective: Maximize expected reward

$$\phi(a) = \sum_{x} R(x) P(x, a)$$
$$P(x, a) = 1 - \prod_{i=1}^{n} (1 - p_i(x, a_i))$$

- Implementation:
 - Assign sensor utilities to induce potential game
 - Apply constrained binary log-linear learning



Self assembly



- References:
 - Yim, Shen, Salemi, Rus, Moll, Lipson, Klavins, & Chirikjian, "Modular self-reconfigurable robot systems: Challenges and opportunities for the future", 2007.
 - Klavins, "Programmable self-assembly", 2007.





- General setup:
 - Nonlocal rules
 - Full "graph grammars"

- Specialized setup:
 - Serial assembly
 - Local rules
 - Bond or break
 - Reversibility



- Complete assembly = Acyclic weighted graph
- Node state: (Position, Vacancies)
- Nodes meet randomly
- If singleton meets vacancy: Active nodes update state
- Singletons break off with probability ϵ

Simulation observation



Critical case: #Atoms = Integer multiple of final assembly

Self assembly & stochastic stability



- Theorem (Fox & JSS, 2009): A state is stochastically stable if and only if there is a minimal number of (sub)assemblies.
- Corollary: Let a complete assembly have N parts. The maximum number of incomplete assemblies is N 1. (For any number of atoms.)

Self assembly proof sketch

- \bullet Form a "backbone" of states with m subassemblies
- Level down: Resistance of 1
- Level up: Resistance at least 2



• Reinforcement learning: x_i = action propensities

$$\begin{aligned} x_i(t+1) &= x_i(t) + \delta(t)(a_i(t) - x_i(t)), \quad \delta(t) = \frac{u_i(a(t))}{t+1} \\ p_i(t) &= (1 - \varepsilon)x_i(t) + \frac{\varepsilon}{N} \mathbf{1} \\ \delta_{\mathsf{std}}(t) &= \frac{u_i(a(t))}{\mathbf{1}^{\mathsf{T}} U_i(t) + u_i(a(t))} \end{aligned}$$

Interpretation: Increased probability of utilized action.

• *Dynamic* reinforcement learning: Introduce running average

$$y_i(t+1) = y_i(t) + \frac{1}{t+1}(x_i(t) - y_i(t))$$
$$p_i(t) = (1-\varepsilon)\Pi_\Delta \left[x_i(t) + \underbrace{\gamma(x_i(t) - y_i(t))}_{\text{new term}} \right] + \frac{\varepsilon}{N} \Pi_\Delta \left[x_i(t) + \underbrace{\gamma(x_i(t) - y_i(t))}_{\text{new term}} \right]$$

• Chasparis & JSS (2009): The pure NE a* has positive probability of convergence iff

$$0 < \gamma_i < \frac{u_i(a_i^*, a_{-i}) - u_i(a_i', a_{-i}^*) + 1}{u_i(a_i', a_{-i}^*)}, \quad \forall a_i' \neq a_i^*$$

(as opposed to all pure NE)

Proof: ODE method of stochastic approximation.

- Implication:
 - Introduction of "forward looking" agent can destabilize equilibria
 - Surviving equilibria = equilibrium selection
- For 2×2 symmetric coordination games
 - RD & not PD \Rightarrow foresight dominance
 - RD & PD & Identical interest \Rightarrow foresight dominance
 - RD & PD together \Rightarrow foresight dominance

- Similar ideas can *stabilize* equilibria (Arslan & JSS)
- Illustration: Perturbed RPS & Marginal foresight replicator dynamics



- Setup:
 - Agents form costly links with other agents
 - Benefits inherited from connectivity

$$u_i(a(t)) = \left(\text{\# of connections to } i \right) - \kappa \cdot \left(\text{\# of links by } i \right)$$

- Properties:
 - Nash networks are "critically connected"
 - Wheel network is unique *efficient* network
 - Chasparis & JSS (2009): The wheel network is foresight dominant.



- Recap:
 - Dynamics & equilibrium selection
 - Prescriptive agenda influence
- Future work:
 - Convergence rates
 - Fully exploit prescriptive agenda
 - Agent dynamics

