Learning Dynamics and Equilibrium Selection

Jeff S Shamma Georgia Institute of Technology

Joint work with George Chasparis, Jason Marden, & Michael Fox

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Networked interaction: Societal, engineered, & hybrid

• How could agents converge to NE?

Arrow: "The attainment of equilibrium requires a disequilibrium process."

- Monographs:
	- Weibull, *Evolutionary Game Theory*, 1997.
	- Young, *Individual Strategy and Social Structure*, 1998.
	- Fudenberg & Levine, *The Theory of Learning in Games*, 1998.
	- Samuelson, *Evolutionary Games and Equilibrium Selection*, 1998.
	- Young, *Strategic Learning and Its Limits*, 2004.
	- Sandholm, *Population Dynamics and Evolutionary Games*, 2010.
- Surveys:
	- Hart, "Adaptive heuristics", *Econometrica*, 2005.
	- Fudenberg & Levine, "Learning and equilibrium", *Annual Review of Economics*, 2009.
- If agents self-organize to Nash equilibrium...
	- Price of Anarchy:

- How to distinguish equlibria?
- Payoff based distinctions: Payoff dominance vs Risk dominance
- Evolutionary (i.e., *dynamic*) distinction
	- Young (1993) "The evolution of convention"
	- Kandori/Mailath/Rob (1993) "Learning, mutation, and long-run equilibria in games"
	- many more...
- Adaptive play:
	- "Two" players sparsely sample from finite history
	- Players either:
		- ∗ Play best response to selection
		- ∗ Experiment with small probability
	- **Young (1993):** Risk dominance is "stochastically stable"
- Dynamics & equilibrium selection theme continued...
	- Constrained log linear learning
	- Self assembly
	- Dynamic reinforcement dynamics
- "Prescriptive" issues & opportunities
	- What are implications of additional constraints?
	- How to exploit additional degrees of freedom?
- Setup:
	- $-$ Players: $\{1, ..., p\}$
	- Actions: $a_i \in \mathcal{A}_i$
	- Action profiles:

$$
(a_1, a_2, ..., a_p) \in \mathcal{A} = \mathcal{A}_1 \times \mathcal{A}_2 \times ... \times \mathcal{A}_p
$$

- Payoffs: $u_i: (a_1,a_2,...,a_p)=(a_i,a_{-i}) \mapsto {\bf R}$
- Global objective: $G : \mathcal{A} \to \mathbf{R}$
- Action profile $a^* \in \mathcal{A}$ is a *Nash equilibrium* (NE) if for all players:

$$
u_i(a_1^*, a_2^*, ..., a_p^*) = u_i(a_i^*, a_{-i}^*) \geq u_i(a_i', a_{-i}^*)
$$

- Learning dynamics:
	- $-t = 0, 1, 2, ...$
	- $\mathbf{Pr} [a_i(t)] = p_i(t), \quad p_i(t) \in \Delta(\mathcal{A}_i)$
	- $-p_i(t) = \mathcal{F}_i$ (available info at time t)

• **Potential games:** For some $\phi : \mathcal{A} \to \mathbb{R}$

$$
\phi(a_i, a_{-i}) - \phi(a'_i, a_{-i}) > 0
$$

$$
\Leftrightarrow
$$

$$
u_i(a_i, a_{-i}) - u_i(a'_i, a_{-i}) > 0
$$

i.e., potential function increases iff unilateral improvement.

- Features:
	- Typical of "coordination games"
	- Desirable convergence properties under various algorithms
	- Need not imply "cooperation" or $\phi = G$
	- Prescriptive opportunity: Potential game by design

Illustrations

- Distributed routing
	- Payoff = negative congestion. $c_r(\sigma_r)$
	- Potential function:

$$
\phi = \sum_{r} \sum_{n=1}^{\sigma_r} c_r(n)
$$

– Overall congestion:

$$
G=\sum_r \sigma_r c_r(\sigma_r)
$$

- $-$ **Note:** $\phi \neq G$
- Multiagent sudoku:

 $u_i(a) = \text{Hreps}$ in row + #reps in column + #reps in sector

$$
\phi(a) = \sum_i u_i(a)
$$

• Preliminary: Gibbs distribution

 $\mathbf{Pr}\left[v_i\right] \propto e^{v_i/T}$

As $T \downarrow 0$ assigns all probability to $\arg \max \{v_1, v_2, ..., v_n\}$

• At stage t

- $-$ Player i is selected at random
- Chosen player selects

$$
\mathbf{Pr}\left[a_i(t) = j\right] \propto e^{u_i(j, a_{-i}(t-1))/T}
$$

– Interpretation: Noisy best reply to previous joint actions

- Fact: SAP results in a Markov chain over joint action space A with a unique stationary distribution, μ .
- **Blume (1993)**: In (cardinal) potential games, steady-state distribution satisfies

$$
\mathbf{Pr}\left[a\right] \propto e^{\phi(a)/T}
$$

- Implication: As $T \downarrow 0$, concentrated at potential maximizer
- i.e., Potential maximizer is "stochastically stable"

• Impose constrained evolution:

$$
a_i(t) \in C(a_i(t-1))
$$

- Limited mobility
- Obstacles
- Mimicking log linear learning *alters* stochastically stable action profiles!
- Example: Identical interest game

L M R U 0 0 9 D 10 -10 -10

$$
C_2(L) = \{L, M\} \quad C_2(M) = \{L, M, R\} \quad C_2(R) = \{M, R\}
$$

- Potential maximizer: (D, L)
- $-$ Stochastically stable state: (U, R)
- Intuition:
	- $*(U, R) \rightarrow (D, L)$ "costs" 18 (used to cost 9)
	- $*(D, L)$ → (U, R) "costs" 10 (used to cost 10)
- At stage t :
	- $-$ Player i is selected at random
	- Chosen player compares $a_i(t-1)$ with randomly selected $a'_i \in C(a_i(t-1))$

 $\mathbf{Pr}\left[a_i(t)\right]\propto e^{u_i(a_i(t-1),a_{-i}(t-1))/T}$ vs $e^{u_i(a_i')}$ t'_{i} , $a_{-i}(t-1)/T$

- **Marden & JSS, 2008:** Under binary log linear learning, only potential function maximizers are stochastically stable.
- *No longer* characterize stationary distribution.
- Recall example:

• Binary version: $(U, M) \rightarrow (U, L)$ now has zero resistance.

- What if evaluation of $U_i(a', a_{-i}(t-1))$ no longer possible?
- New setup: Players can only measure $a_i(t)$ and $U_i(a(t))$
- \bullet Introduce *baseline action* a_i^b $\boldsymbol{h}_i^b(t)$ and *baseline utility* \boldsymbol{u}_i^b $i^b(t)$
- Action selection:

$$
a_i(t) = a_i^b(t)
$$
 with probability $(1 - \epsilon)$

 $a_i(t)$ is chosen randomly over \mathcal{A}_i with probability ϵ

• Baseline action & utility update:

New baseline with probability $\sim e^{U_i(a(t))/T}$ ⇓ a_i^b $i(t+1) = a_i(t)$ u_i^b $u_i^b(t+1) = u_i(a(t))$ u_i^b *Keep baseline* with probability $\sim e^{u^b_i}$ $\frac{b}{i}(t)/T$ ⇓ a_i^b $a_i^b(t+1) = a_i^b$ $\frac{b}{i}(t)$ $i(t+1) = u_i^b$ $i^b(t)$

• **Marden & JSS, 2008:** Under payoff based linear learning, only potential function maximizers are stochastically stable.

- Definition:
	- Let P^{ϵ} denote the transition probability matrix of an irreducible & aperiodic Markov chain.
	- Let μ^{ϵ} be the (unique) stationary distribution for P^{ϵ}
	- A state, x, is **stochastically stable** if

$$
\liminf_{\epsilon \to 0} \mu^{\epsilon}(x) > 0
$$

- **Young (1993)**: To determine stochastic stability
	- View learning dynamics as ϵ perturbation of reference ($\epsilon = 0$) Markov chain
	- Divide reference Markov chain into recurrence classes
	- Define *resistance* to transition between recurrence classes:

$$
0<\lim_{\epsilon\downarrow 0}\frac{P_{ij}^\epsilon}{\epsilon^{r(i\to j)}}<\infty
$$

- Form *stochastic potential* for each recurrence class
- Minimal stochastic potential implies stochastic stability

• Combinatoric utilization vs pragmatic utilization

Illustration: Sensor allocation

• Objective: Maximize expected reward

$$
\phi(a) = \sum_{x} R(x)P(x, a)
$$

$$
P(x, a) = 1 - \Pi_{i=1}^{n} (1 - p_i(x, a_i))
$$

- Implementation:
	- Assign sensor utilities to induce potential game
	- Apply constrained binary log-linear learning

Self assembly

- References:
	- Yim, Shen, Salemi, Rus, Moll, Lipson, Klavins, & Chirikjian, "Modular self-reconfigurable robot systems: Challenges and opportunities for the future", 2007.
	- Klavins, "Programmable self-assembly", 2007.

- General setup:
	- Nonlocal rules
	- Full "graph grammars"
- Specialized setup:
	- Serial assembly
	- Local rules
	- Bond or break
	- Reversibility

Assembly rules

- Complete assembly = Acyclic weighted graph
- Node state: (Position, Vacancies)
- Nodes meet randomly
- If singleton meets vacancy: Active nodes update state
- Singletons break off with probability ϵ

Simulation observation

Critical case: #Atoms = Integer multiple of final assembly

Self assembly & stochastic stability

- **Theorem (Fox & JSS, 2009):** A state is stochastically stable if and only if there is a minimal number of (sub)assemblies.
- Corollary: Let a complete assembly have N parts. The maximum number of incomplete assemblies is $N - 1$. (For any number of atoms.)

Self assembly proof sketch

- Form a "backbone" of states with m subassemblies
- Level down: Resistance of 1
- Level up: Resistance at least 2

• Reinforcement learning: x_i = action propensities

$$
x_i(t+1) = x_i(t) + \delta(t)(a_i(t) - x_i(t)), \quad \delta(t) = \frac{u_i(a(t))}{t+1}
$$

$$
p_i(t) = (1 - \varepsilon)x_i(t) + \frac{\varepsilon}{N}\mathbf{1}
$$

$$
\delta_{\text{std}}(t) = \frac{u_i(a(t))}{\mathbf{1}^{\text{T}}U_i(t) + u_i(a(t))}
$$

Interpretation: Increased probability of utilized action.

• *Dynamic* reinforcement learning: Introduce running average

$$
y_i(t+1) = y_i(t) + \frac{1}{t+1}(x_i(t) - y_i(t))
$$

$$
p_i(t) = (1 - \varepsilon)\Pi_\Delta \left[x_i(t) + \underbrace{\gamma(x_i(t) - y_i(t))}_{\text{new term}} \right] + \frac{\varepsilon}{N} \mathbf{1}
$$

• **Chasparis & JSS (2009):** The pure NE a^* has positive probability of convergence iff

$$
0 < \gamma_i < \frac{u_i(a_i^*, a_{-i}) - u_i(a_i', a_{-i}^*) + 1}{u_i(a_i', a_{-i}^*)}, \quad \forall a_i' \neq a_i^*
$$

(as opposed to all pure NE) *Proof: ODE method of stochastic approximation.*

- Implication:
	- Introduction of "forward looking" agent can destabilize equilibria
	- Surviving equilibria = equilibrium selection
- For 2×2 symmetric coordination games
	- RD & not $PD \Rightarrow$ foresight dominance
	- RD & PD & Identical interest \Rightarrow foresight dominance
	- RD & PD together \neq foresight dominance
- Similar ideas can *stabilize* equilibria (Arslan & JSS)
- Illustration: Perturbed RPS & Marginal foresight replicator dynamics

- Setup:
	- Agents form costly links with other agents
	- Benefits inherited from connectivity

$$
u_i(a(t)) = \Bigl(\# \text{ of connections to } i\Bigr) - \kappa \cdot \Bigl(\# \text{ of links by } i\Bigr)
$$

- Properties:
	- Nash networks are "critically connected"
	- Wheel network is unique *efficient* network
	- **Chasparis & JSS (2009):** The wheel network is foresight dominant.

- Recap:
	- Dynamics & equilibrium selection
	- Prescriptive agenda influence
- Future work:
	- Convergence rates
	- Fully exploit prescriptive agenda
	- Agent dynamics

