



# Controllability metrics, limitations and algorithms for complex networks

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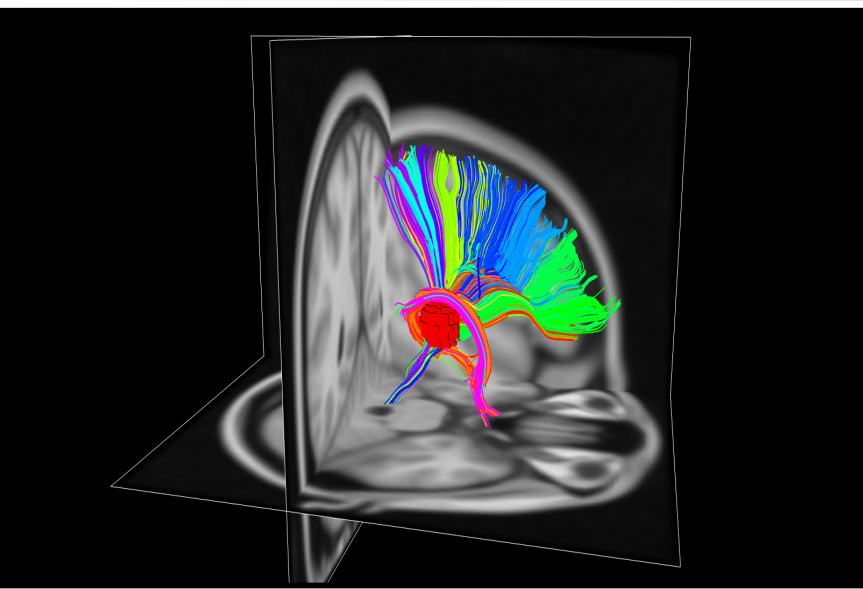
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In collaboration with

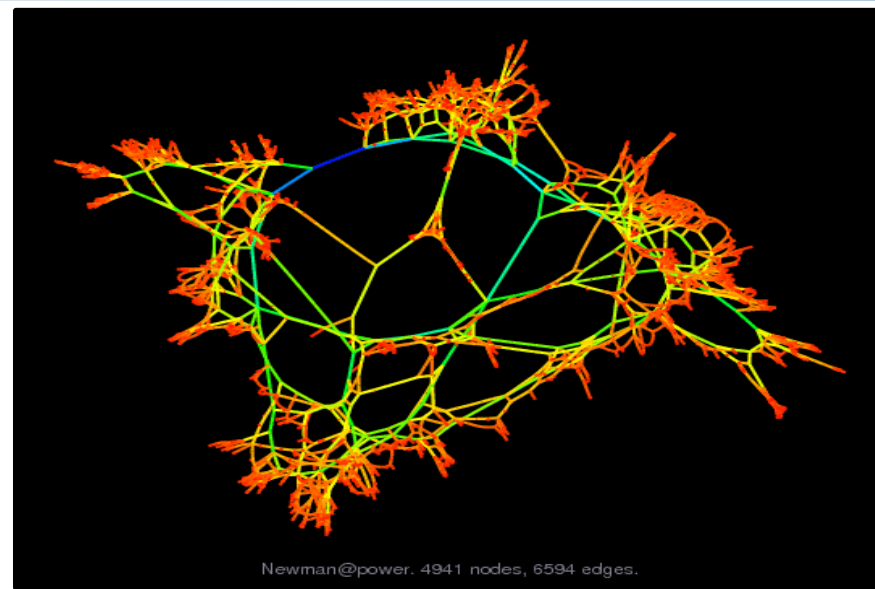
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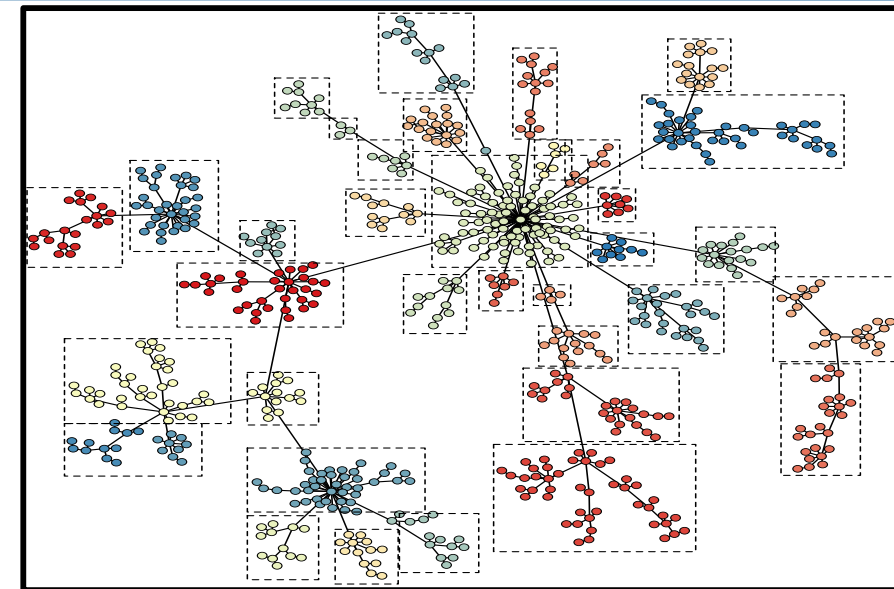
# Growing Number of Applications



Brain network



Power network



Social network

- ✓ Y. Y. Liu, J. J. Slotine, and A. L. Barabasi, “Controllability of complex networks,” Nature 2011.
- ✓ N. J. Cowan, E. J. Chastain, D. A. Vilhena, J. S. Freudenberg, and C. T. Bergstrom, “Nodal dynamics, not degree distributions, determine the structural controllability of complex networks,” PLOS ONE, 2012.
- ✓ G. Yan, J. Ren, Y.-C. Lai, C.-H. Lai, and B. Li, “Controlling complex networks: How much energy is needed?” Physical Review Letters, 2012.
- ✓ J. Sun and A. E. Motter, “Controllability transition and nonlocality in network control,” Physical Review Letters, 2013.
- ✓ F. L. Cortesi, T. H. Summers, and J. Lygeros, “Submodularity of Energy Related Controllability Metrics,” Arxiv preprint, 2014.
- ✓ A. Olshevsky, “Minimal Controllability Problems,” Arxiv preprint, 2014.

# Problem formulation

Consider a linear system

$$x(t + 1) = Ax(t) + Bu(t)$$

where  $A$  is sparse  $n \times n$  matrix (the interactions between the states is described by a graph) and

$$B = [e_{i_1} \cdots e_{i_m}]$$

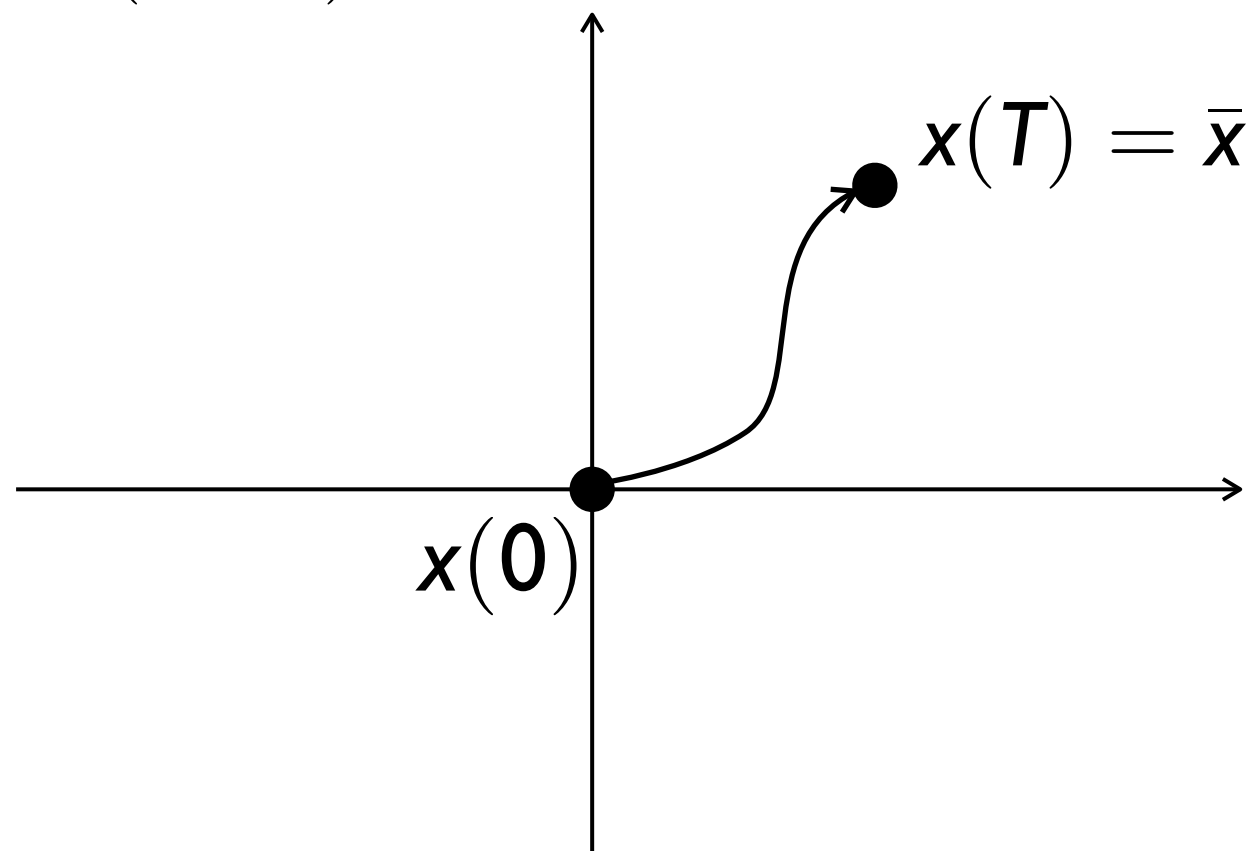
where

$$e_i = \begin{bmatrix} 0 \\ \vdots \\ 0 \\ 1 \\ 0 \\ \vdots \\ 0 \end{bmatrix} \leftarrow i$$

# Problem formulation

$$x(t + 1) = Ax(t) + Bu(t)$$

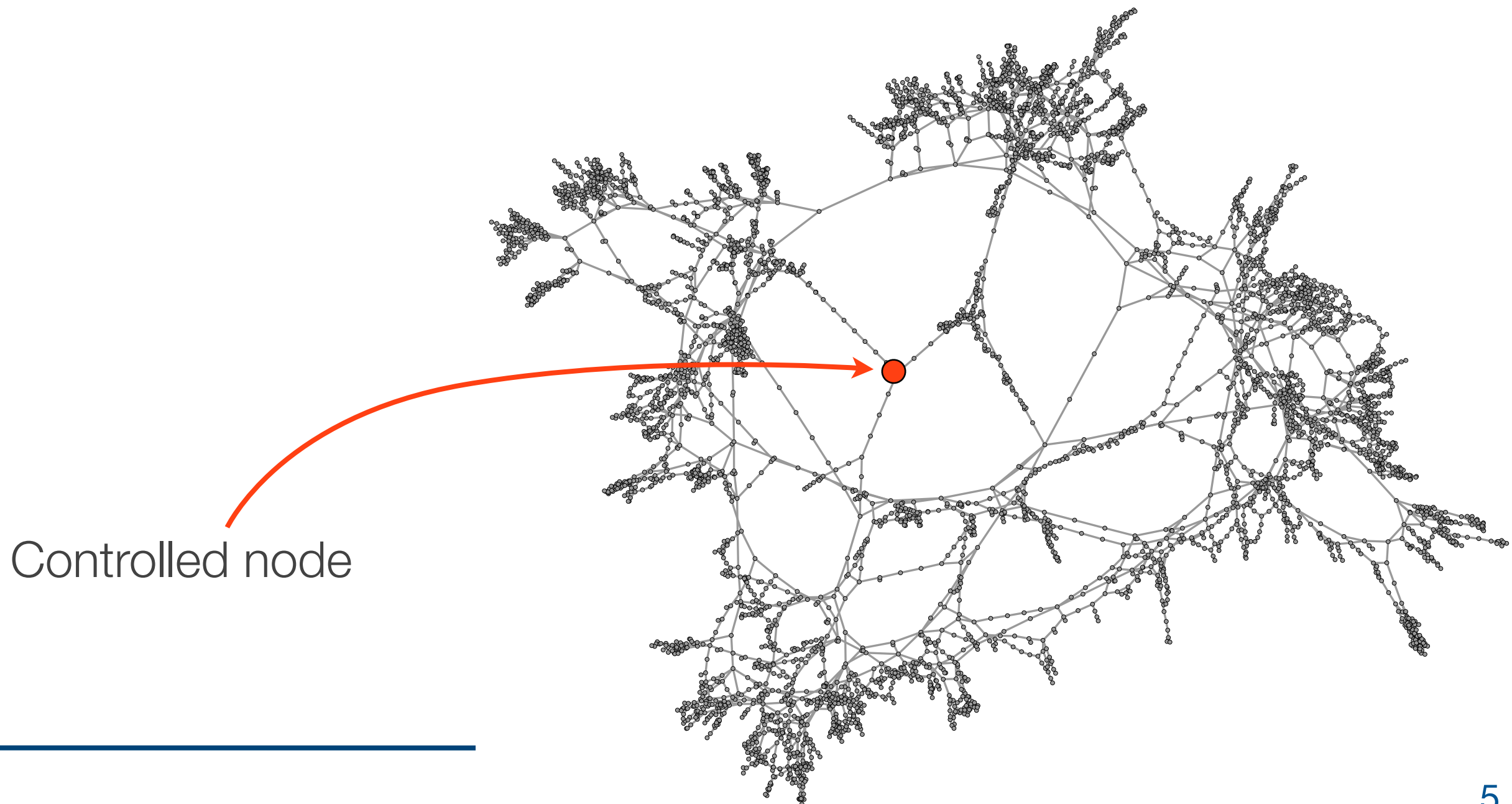
**CONTROLLABILITY** is the possibility of steering the state from the initial state  $x(0) = 0$  to an arbitrary final state  $x(T) = \bar{x}$  by applying a suitable input sequence  $u(0), u(1), \dots, u(T - 1)$ .





# Need of a controllability metric

Assume that the graph is strongly connected and that it has all the self-loops. Then to have that the resulting system is controllable (generically in the non-zero entries of  $A$ ) it is enough that only one state is controlled.





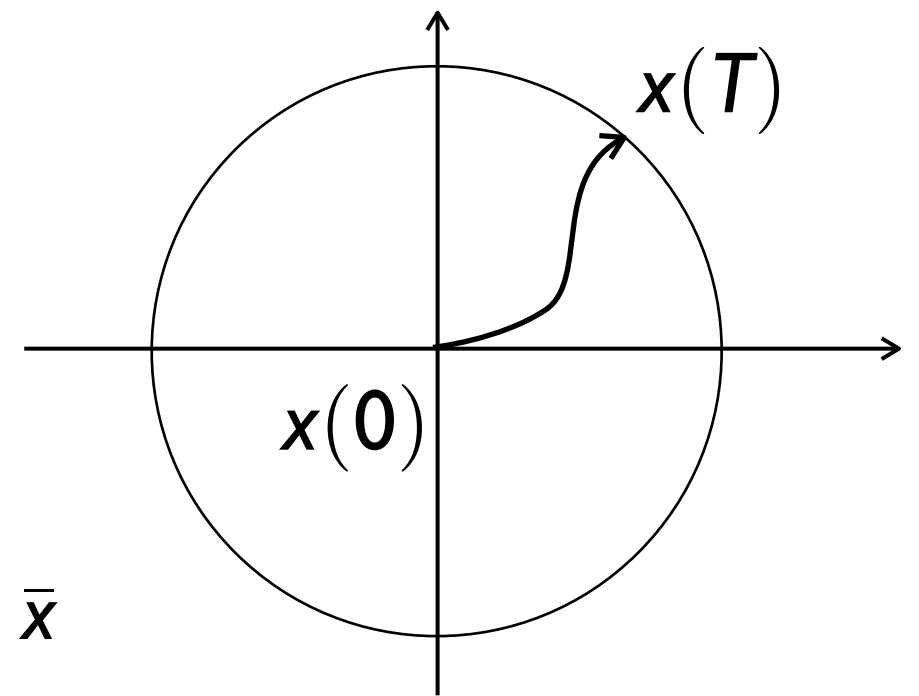
# Need of a controllability metric

**QUESTION:** How controllable is the resulting system?

# A controllability metric

One metric for describing the controllability degree is given by the controllability Gramian

$$W_T := \sum_{t=0}^{T-1} A^t B B^T (A^T)^t$$



**Interpretation:** For driving

$$x(0) = 0 \quad \longrightarrow \quad x(T) = \bar{x}$$

where  $\|\bar{x}\|_2 = 1$ , we need an input  $u(0), u(1), \dots, u(T-1)$  with  $L^2$  energy  $\bar{x}^T W_T^{-1} \bar{x}$ .

The energy to drive the state from zero to a norm one state (in the worst case) is given by

$$\text{Energy} = \frac{1}{\lambda_{\min}(W_T)}$$

# A controllability metric

$$\mathbf{W}_T := \sum_{t=0}^{T-1} \mathbf{A}^t \mathbf{B} \mathbf{B}^T (\mathbf{A}^T)^t \quad \text{Controllability Gramian}$$

$$\text{Energy} = \frac{1}{\lambda_{\min}(\mathbf{W}_T)}$$

Small  $\lambda_{\min}(\mathbf{W}_T)$   $\longleftrightarrow$  Small controllability degree

Large  $\lambda_{\min}(\mathbf{W}_T)$   $\longleftrightarrow$  Large controllability degree



# A controllability metric

## Alternative controllability metrics

$$\mathbf{W}_T := \sum_{t=0}^{T-1} \mathbf{A}^t \mathbf{B} \mathbf{B}^T (\mathbf{A}^T)^t$$

$$\lambda_{\min}(\mathbf{W}_T)$$

$$\frac{1}{n} \text{tr}(\mathbf{W}_T^{-1})$$

$$\frac{1}{n} \ln \det(\mathbf{W}_T)$$

$$\frac{1}{n} \text{tr}(\mathbf{W}_T)$$

# Few Nodes Cannot Control Symmetric Complex Networks

## THEOREM

Assume the matrix  $A$  symmetric. Fix any constant  $0 < \mu < 1$  and let

$$n(\mu) := |\{\lambda \in \lambda(A) : |\lambda|^2 \leq \mu\}|$$

Then

$$\lambda_{\min}(\mathbf{W}_T) \leq \frac{1}{\mu(1-\mu)} \mu^{n(\mu)/m}$$

# Few Nodes Cannot Control Symmetric Complex Networks

Let

$$F_A(\bar{\lambda}) := \frac{|\{\lambda \in \lambda(\mathbf{A}) : \lambda \leq \bar{\lambda}\}|}{n}$$

and

$$f_A(\bar{\lambda}) := \frac{dF(\bar{\lambda})}{d\bar{\lambda}}$$

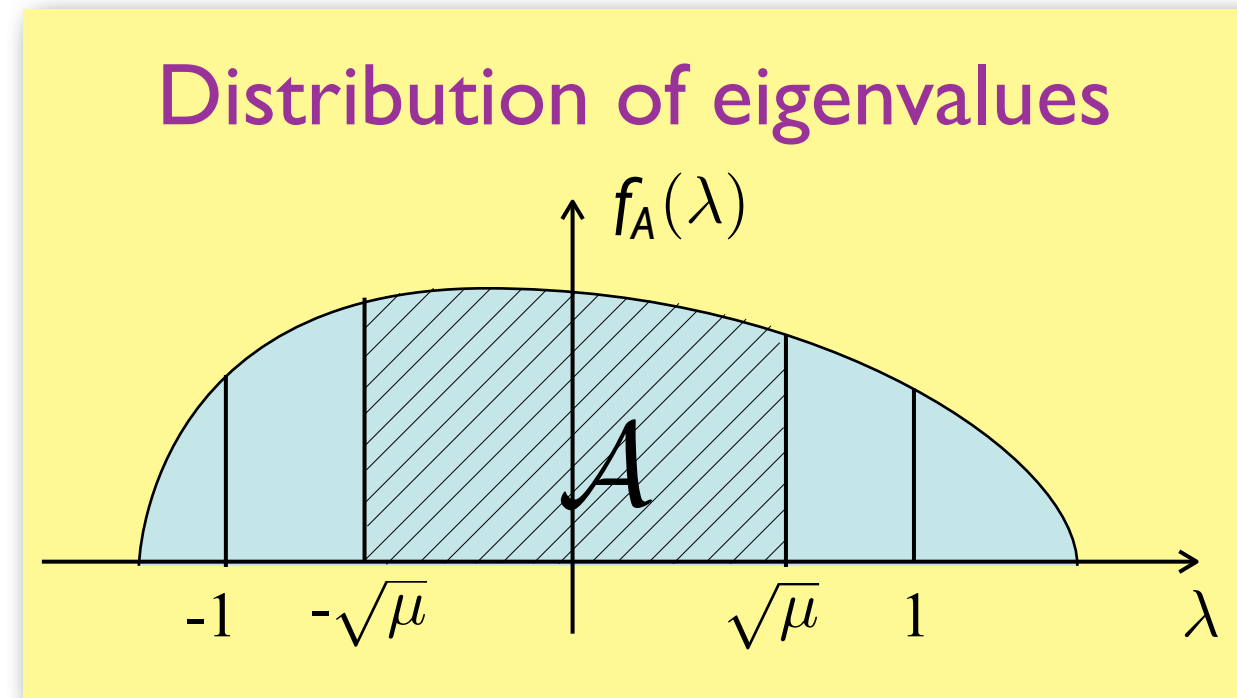
Then

$$n(\mu) := |\{\lambda \in \lambda(\mathbf{A}) : |\lambda|^2 \leq \mu\}| = n \int_{-\sqrt{\mu}}^{\sqrt{\mu}} f_A(\lambda) d\lambda = \mathcal{A} n$$

Hence the number  $n(\mu)$  typically grows linearly in the network cardinality  $n$ .

$$\lambda_{\min}(\mathbf{W}_T) \leq \frac{1}{\mu(1-\mu)} \mu^{\mathcal{A} \frac{n}{m}}$$

$$\mu < 1$$

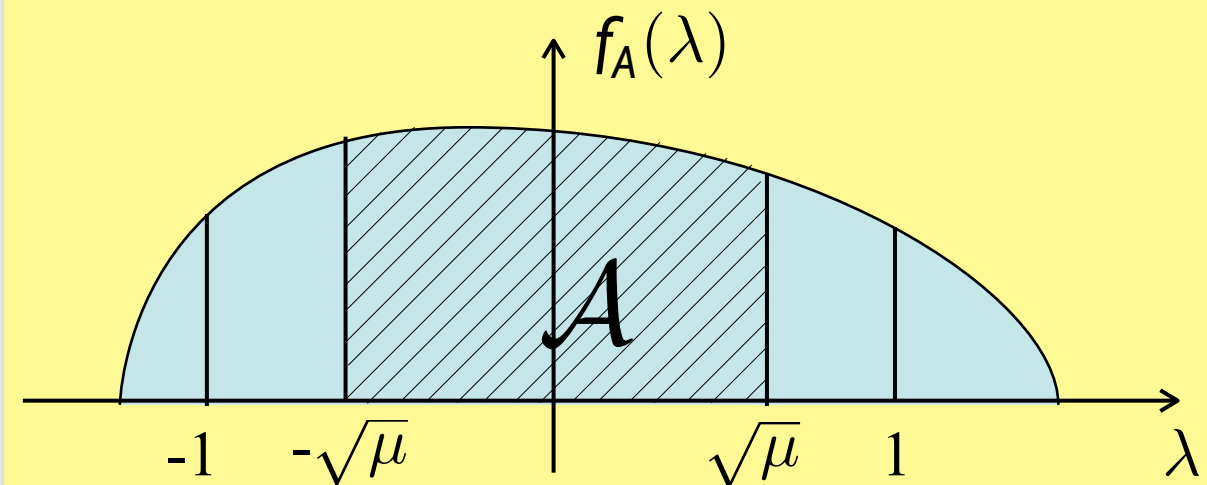


# Few Nodes Cannot Control Symmetric Complex Networks

$$\mu < 1$$

$$\lambda_{\min}(\mathbf{W}_T) \leq \frac{1}{\mu(1-\mu)} \mu^A \frac{n}{m}$$

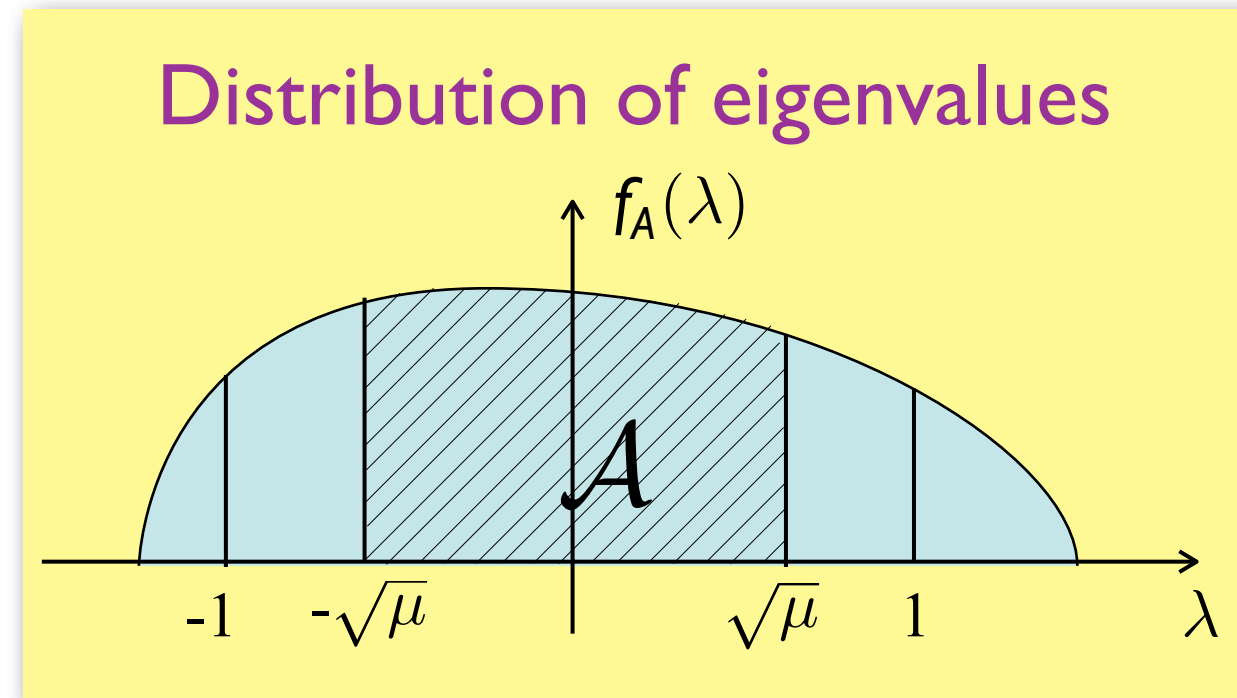
Distribution of eigenvalues



# Few Nodes Cannot Control Symmetric Complex Networks

$$\mu < 1$$

$$\lambda_{\min}(\mathbf{W}_T) \leq \frac{1}{\mu(1-\mu)} \mu^A \left(\frac{n}{m}\right)$$

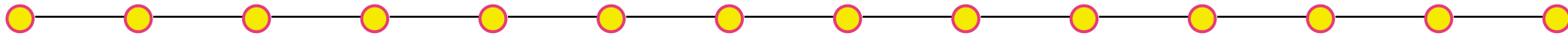


✓ For fixed number  $m$  of control nodes, the controllability degree **decreases exponentially** in the network cardinality  $n$ .

✓ To have a fixed controllability degree, number  $m$  of control nodes must grow **linearly** in the network cardinality  $n$ .



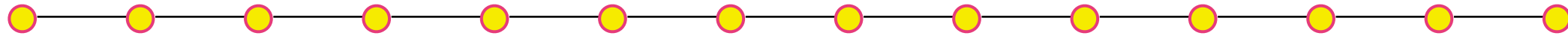
# Example: symmetric line graph



$$A = \begin{bmatrix} a & b & 0 & \dots & \dots & 0 & 0 \\ b & a & b & \dots & \dots & 0 & 0 \\ 0 & b & a & \dots & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \ddots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \ddots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & \dots & a & b \\ 0 & 0 & 0 & \dots & \dots & b & a \end{bmatrix}$$

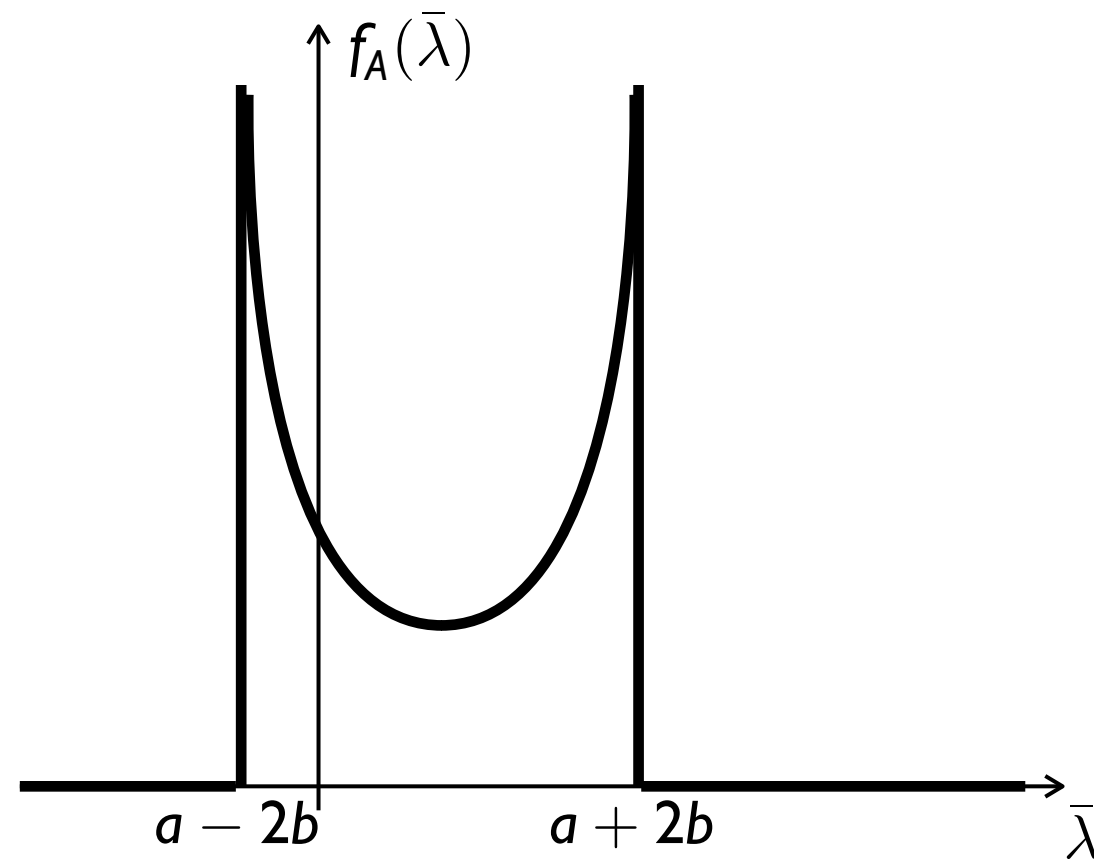
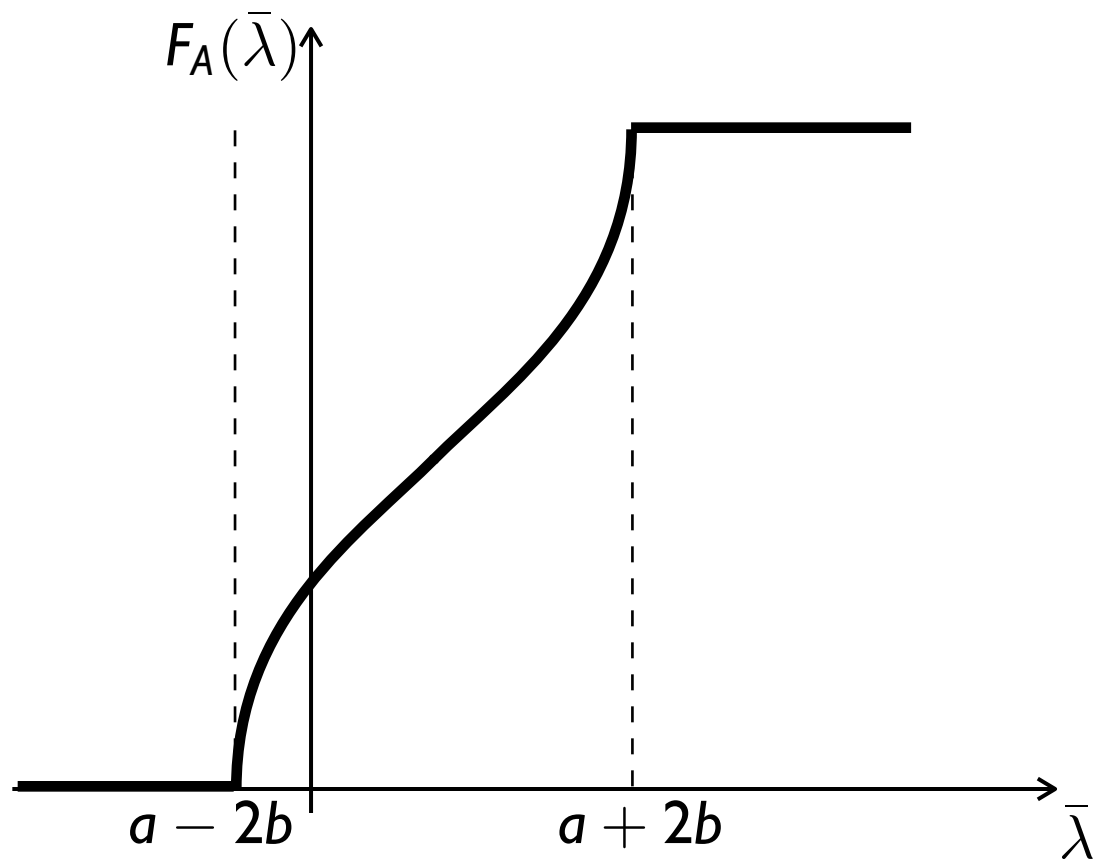
$$\lambda(A) = \left\{ a + 2b \cos \left( \frac{\pi}{n+1} k \right) : k = 1, 2, \dots, n \right\}$$

# Example: symmetric line graph

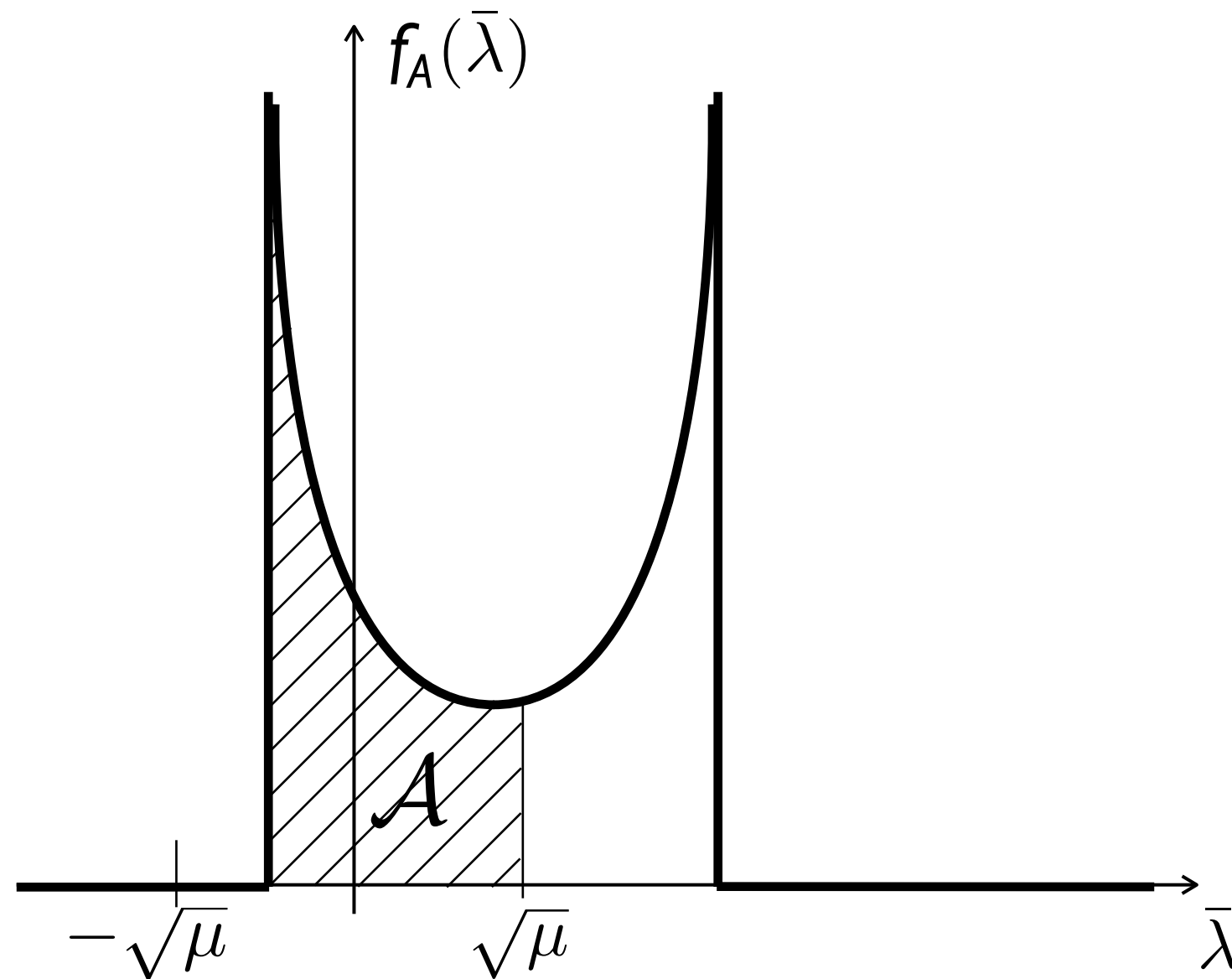
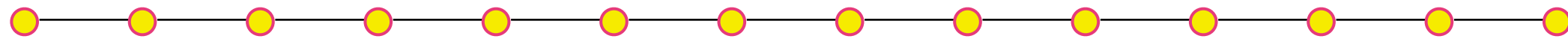


$$F_A(\bar{\lambda}) = 1 - \frac{1}{\pi} \arccos \left( \frac{\bar{\lambda} - a}{2b} \right)$$

$$f_A(\bar{\lambda}) = \frac{1/\pi}{\sqrt{1 - \left( \frac{\bar{\lambda} - a}{2b} \right)^2}}$$



# Example: symmetric line graph



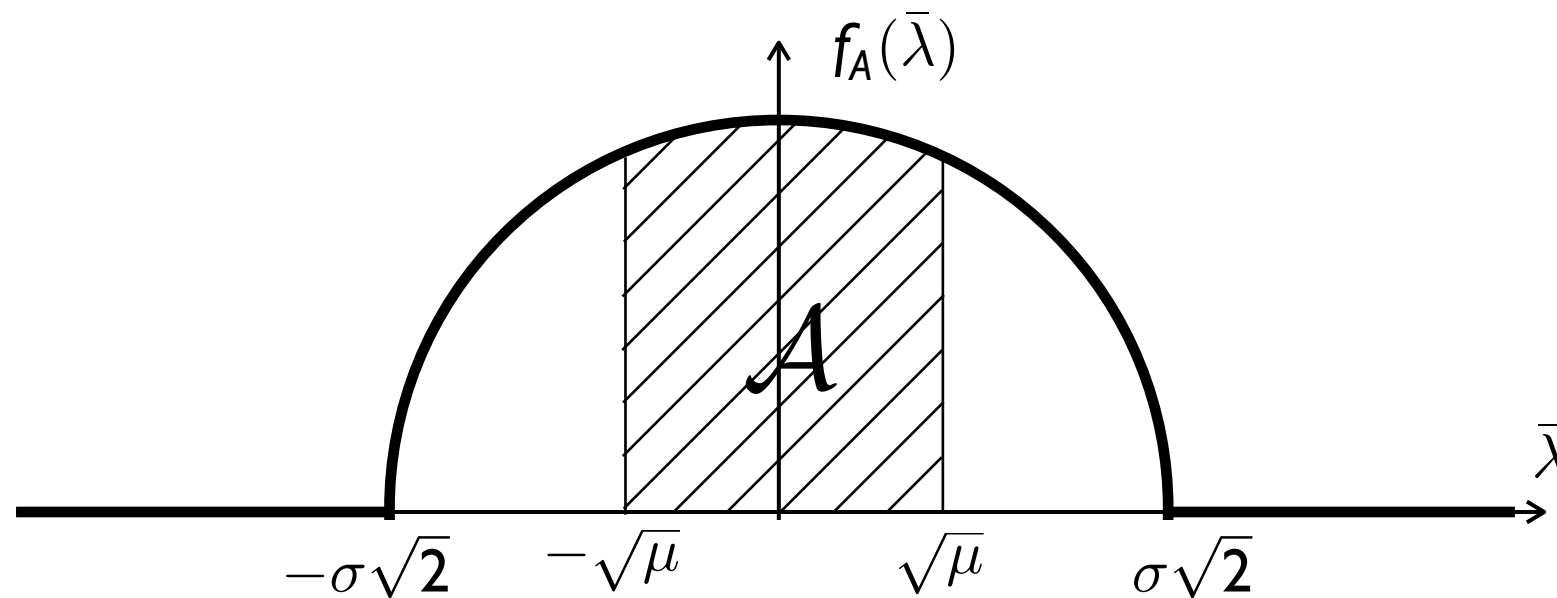
# Example: Symmetric random matrix

Assume  $A$  is a symmetric matrix with  $a_{ij}$  iid random variables  $\mathbb{E}[a_{ij}] = 0$  and  $\mathbb{E}[a_{ij}^2] = \sigma^2/\sqrt{n}$ .

Then (Wigner's semi-circle law)

$$f_A(\bar{\lambda}) \longrightarrow W(\bar{\lambda})$$

$$A = \begin{bmatrix} a_{11} & a_{12} & \cdots & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & \cdots & a_{nn} \end{bmatrix}$$





# Asymmetric Complex Networks

For asymmetric networks the situation is more complex

- For “**isotropic**” networks (networks with no preferential directions) it seems that the situation is the same as for symmetric networks, namely they are difficult to control.
- For “**anisotropic**” networks (networks with a preferential direction) it seems that few nodes can indeed control large scale networks.



## THEOREM

Assume the matrix  $A$  diagonalizable and let  $V$  an eigenvector matrix. Fix any constant  $0 < \mu < 1$  and let

$$n(\mu) := |\{\lambda \in \lambda(A) : |\lambda|^2 \leq \mu\}|$$

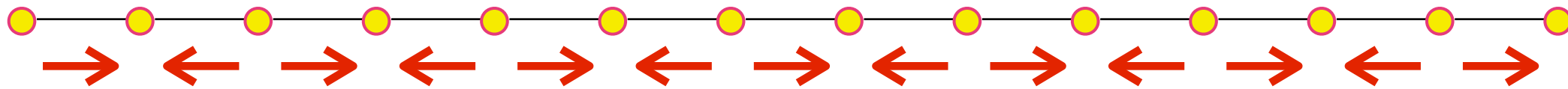
Then

$$\lambda_{\min}(W_T) \leq \text{cond}(V)^2 \frac{1}{\mu(1-\mu)} \mu^{n(\mu)/m}$$

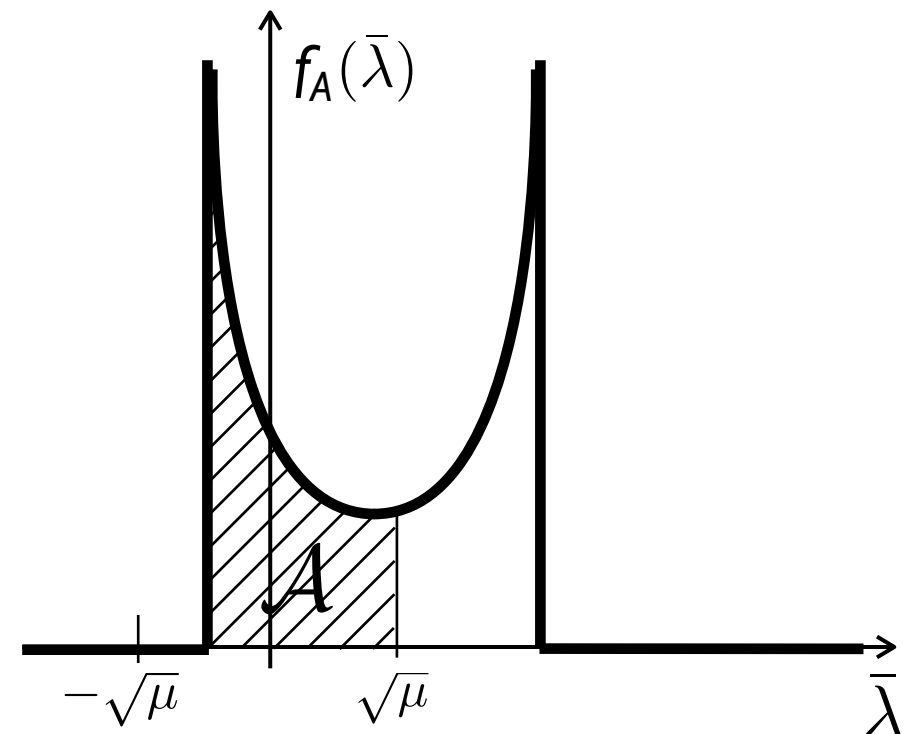
where  $\text{cond}(V) = \sigma_{\max}(V)/\sigma_{\min}(V)$  is the condition number of  $V$ .

**Consequence:** If the condition number stays bounded in the network dimension, than the network remains difficult to control.

# Example: Asymmetric isotropic line graph



$$A := \frac{1}{3} \begin{bmatrix} 1 & 1/2 & 0 & 0 & \dots \\ 2 & 1 & 2 & 0 & \dots \\ 0 & 1/2 & 1 & 1/2 & \dots \\ 0 & 0 & 2 & 1 & \dots \\ \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}.$$

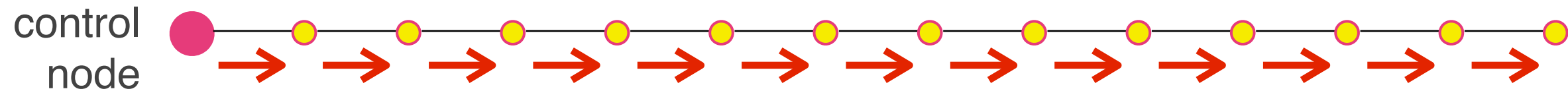


It can be found an eigenvector matrix  $V$  such that  $\text{cond}(V) = 2$ .

$$\lambda(A) = \left\{ \frac{1}{3} + \frac{2}{3} \cos \frac{k\pi}{n+1}, \quad k = 1, \dots, n \right\}$$

This network is difficult to control as a symmetric network.

# Example: Asymmetric anisotropic line graph



$$A = \begin{bmatrix} a & b & 0 & \dots & \dots & 0 & 0 \\ c & a & b & \dots & \dots & 0 & 0 \\ 0 & c & a & \dots & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \ddots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \ddots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & \dots & a & b \\ 0 & 0 & 0 & \dots & \dots & c & a \end{bmatrix}$$

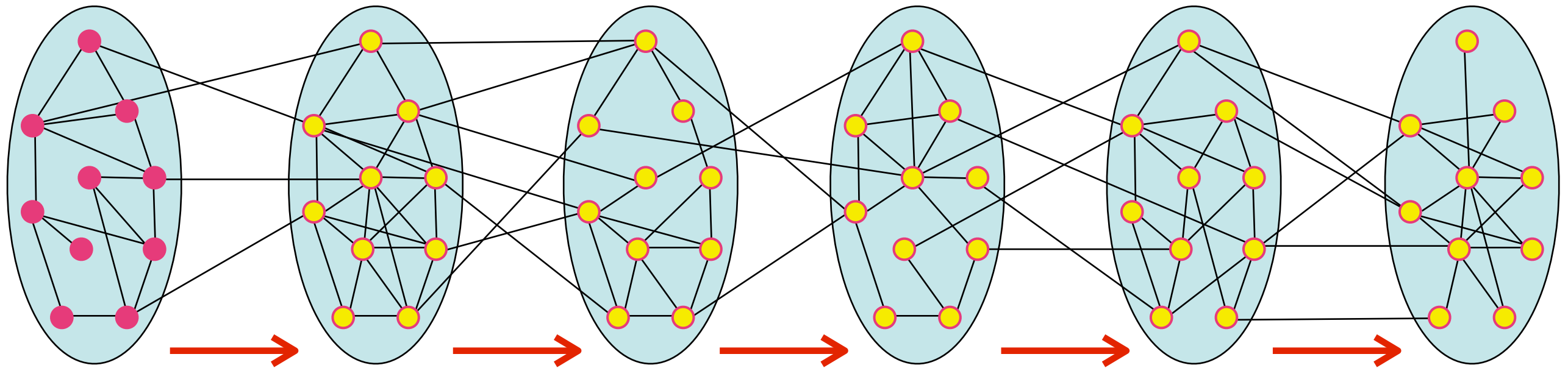
$$\lambda(A) = \left\{ a + 2\sqrt{bc} \cos \left( \frac{\pi}{n+1} k \right) : k = 1, 2, \dots, n \right\}$$

Given  $a$ , if  $c$  is sufficiently larger than  $b$ , then this network can be controlled with finite energy by the node on the extreme regardless the network dimension.

**Exploiting spatial instability**

# Extension to more general graphs: controllable graphs

control nodes

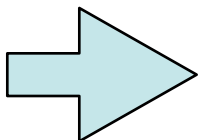


$$A = \begin{bmatrix} A_1 & B_1 & 0 & \dots & \dots & 0 & 0 \\ C_1 & A_2 & B_2 & \dots & \dots & 0 & 0 \\ 0 & C_2 & A_3 & \dots & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \ddots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \ddots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & \dots & A_{l-1} & B_{l-1} \\ 0 & 0 & 0 & \dots & \dots & C_{l-1} & A_l \end{bmatrix}$$

# Extension to more general graphs

$$A = \begin{bmatrix} A_1 & B_1 & 0 & \dots & \dots & 0 & 0 \\ C_1 & A_2 & B_2 & \dots & \dots & 0 & 0 \\ 0 & C_2 & A_3 & \dots & \dots & 0 & 0 \\ \vdots & \vdots & \vdots & \ddots & \ddots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \ddots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & \dots & \dots & A_{l-1} & B_{l-1} \\ 0 & 0 & 0 & \dots & \dots & C_{l-1} & A_l \end{bmatrix}$$

Given the matrices  $A_i$ , if the matrices  $C_i$  are sufficiently larger than the matrices  $B_i$ , then this network can be controlled with finite energy by the nodes on the extreme subgraph regardless the network dimension.





## Extension to more general graphs: uncontrollable graphs

Assume  $A$  is a stochastic matrix (diffusion dynamics, consensus, ... ). This means that

$$A\mathbf{1} = \mathbf{1}$$

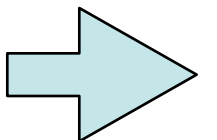
where  $\mathbf{1}$  is the vector with entries equal to one. Let  $w$  be a vector such that

$$w^T A = w^T \quad w^T \mathbf{1} = 1$$

which means that  $w$  is the invariant measure of the Markov chain associated with  $A$ . It is known that the entries of  $w$  represent the nodes "centrality" in the network (the bigger the more important).

**Result:** If the entries of  $w$  are all  $\leq \frac{\text{cost}}{n}$  (all nodes have similar centrality) then the associated network is difficult to control.

In the symmetric case the entries of  $w$  are  $1/n$ .



# Controllers positioning

## Decoupled control strategy: "Divide et Impera"

**Network partitioning** Partition  $\mathcal{V} = \{1, \dots, n\}$  into  $N$  disjoint sets  $\mathcal{V}_1, \dots, \mathcal{V}_N$ . After relabeling of states and inputs, the matrices read as

$$A = \begin{bmatrix} A_1 & \cdots & A_{1N} \\ \vdots & \vdots & \vdots \\ A_{N1} & \cdots & A_N \end{bmatrix}, \quad B = \begin{bmatrix} B_1 & \cdots & 0 \\ \vdots & \ddots & \vdots \\ 0 & \cdots & B_N \end{bmatrix},$$

The networks dynamics can be written as the interconnection of  $N$  subsystems of the form

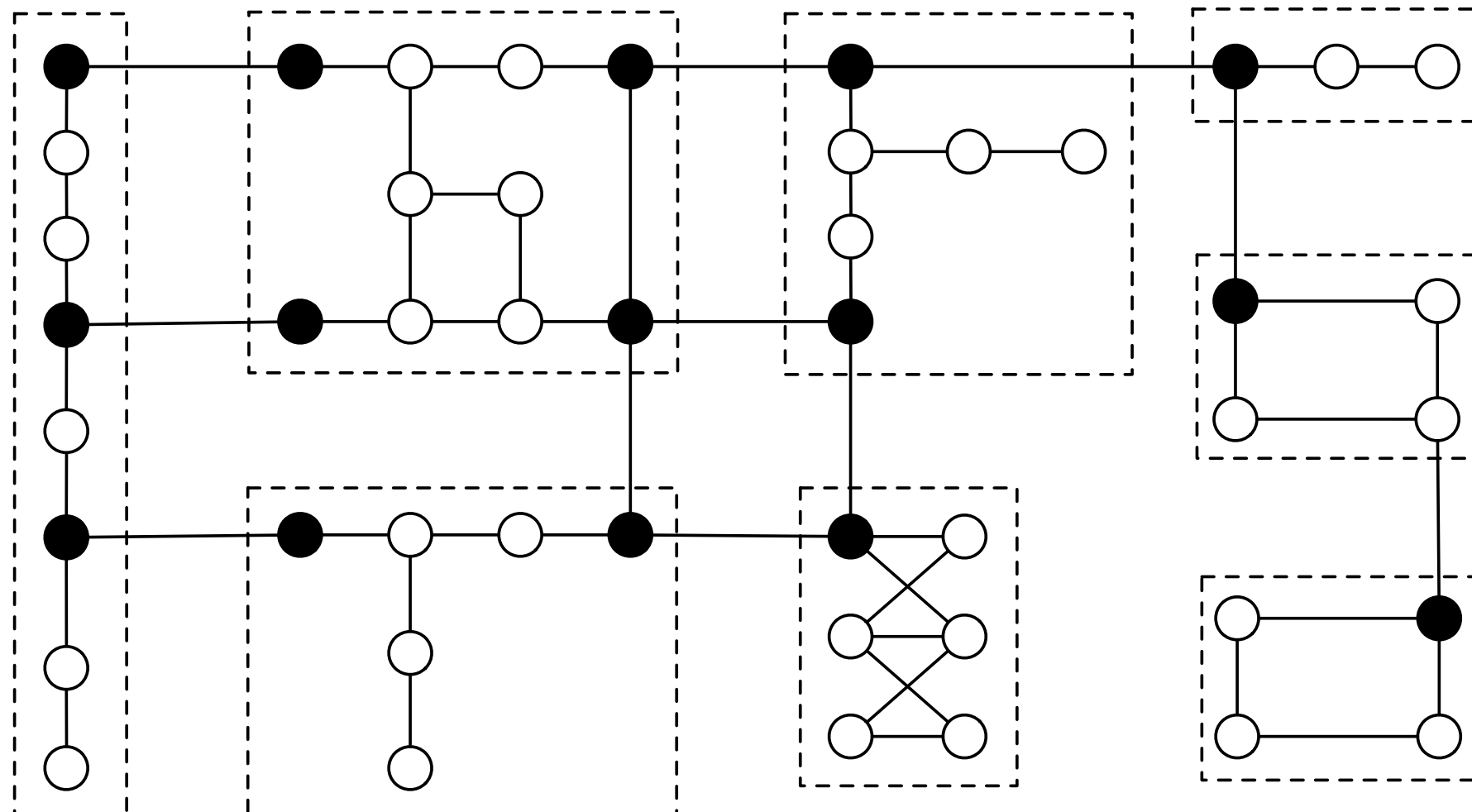
$$x_i(t+1) = \overbrace{A_i x_i(t)}^{\text{local dynamics}} + \overbrace{\sum_{j \in \mathcal{N}_i} A_{ij} x_j(t)}^{\text{interconnections}} + \overbrace{B_i u_i(t)}^{\text{local controls}},$$

where  $i \in \{1, \dots, N\}$  and  $\mathcal{N}_i := \{j : A_{ij} \neq 0\}$ .

# Controllers positioning

## Decoupled control strategy

1. Partition the network into disjoint connected parts.
2. Select boundary nodes as control nodes.



## Decoupled control strategy

### 3. Apply the inputs

$$u_i(t) := v_i(t) - \sum_{j \in \mathcal{N}_i} B_i^T A_{ij} x_j(t)$$

This control law yields  $N$  decoupled subsystems

$$x_i(t + 1) = A_i x_i(t) + B_i v_i(t)$$

4. Choose  $v_i$  which minimizes the energy to steer the subsystem to the desired substate.

# Controllers positioning

## Local controllability

$$\Lambda := \begin{bmatrix} \lambda_{\min}^{-1}(W_{1,T}) & \mathbf{0}_2 & \cdots & \mathbf{0} \\ \mathbf{0} & \lambda_{\min}^{-1}(W_{2,T}) & \cdots & \mathbf{0} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{0} & \mathbf{0} & \cdots & \lambda_{\min}^{-1}(W_{N,T}) \end{bmatrix}.$$

## Coupling strength

$$\Delta := \begin{bmatrix} \mathbf{I} & \|A_{12}\|_2 & \cdots & \|A_{1N}\|_2 \\ \|A_{21}\|_2 & \mathbf{I} & \cdots & \|A_{2N}\|_2 \\ \vdots & \vdots & \ddots & \vdots \\ \|A_{N1}\|_2 & \|A_{N2}\|_2 & \cdots & \mathbf{I} \end{bmatrix}.$$



# Controllers positioning

**Theorem** If we choose a decoupled control law then we obtain

$$\lambda_{\min}(\mathbf{W}_T) \geq \frac{(1 - \bar{\lambda}_{\max})^2}{\|\Lambda\|_{\infty} \|\Delta\|_{\infty}^2}$$

where

$$\bar{\lambda}_{\max} = \max\{\lambda_{\max}(A_i) : i \in \{1, \dots, N\}\} < 1$$

# Controllers positioning

**Theorem** If we choose a decoupled control law  
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where

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# Controllers positioning

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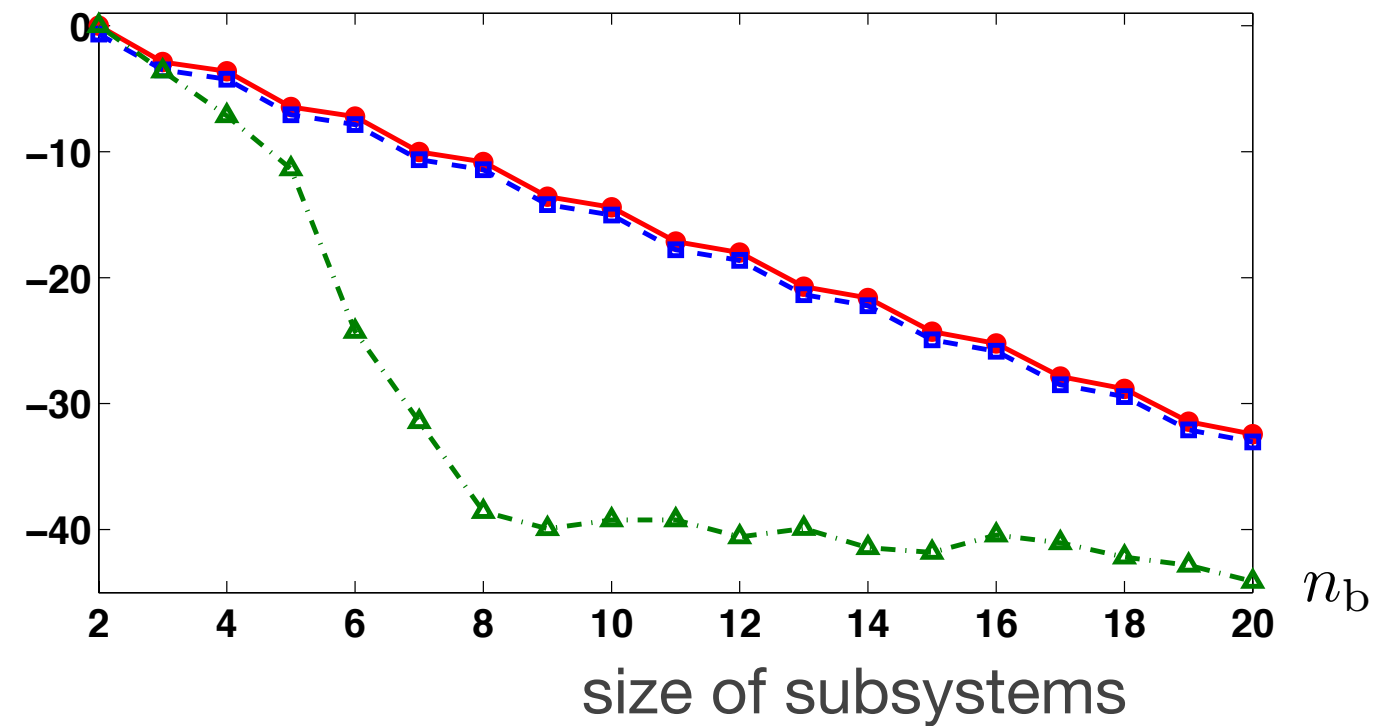
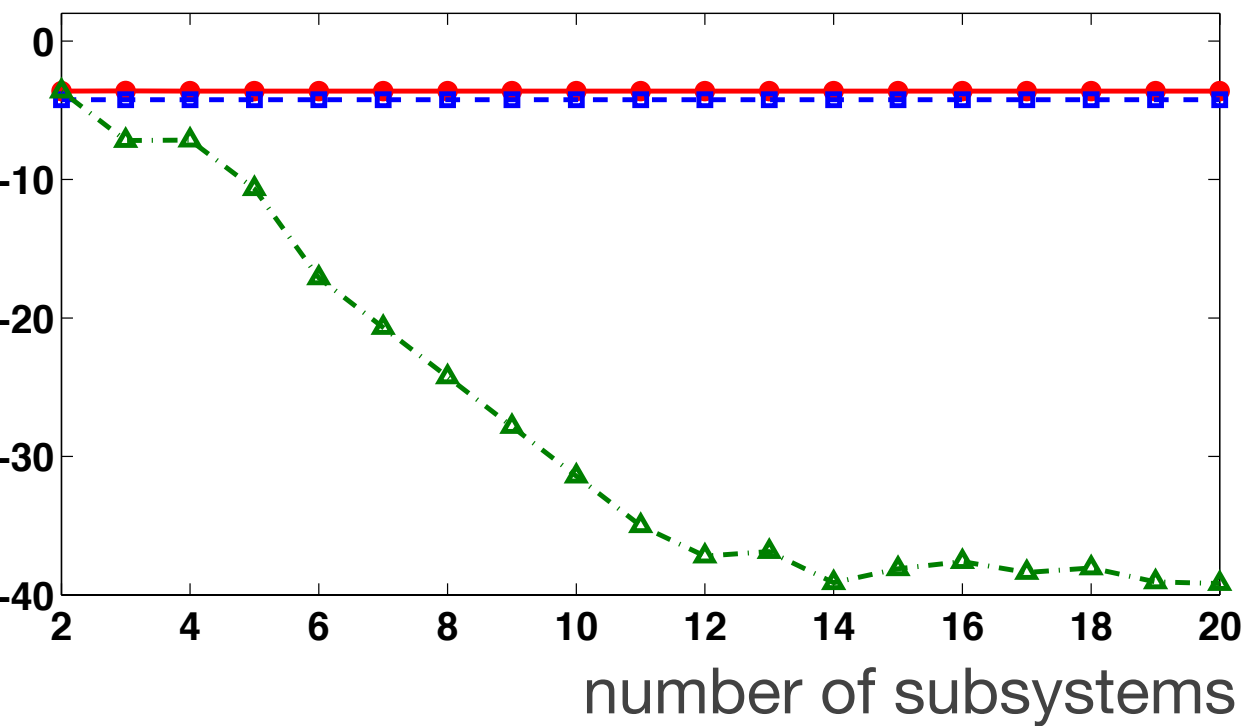
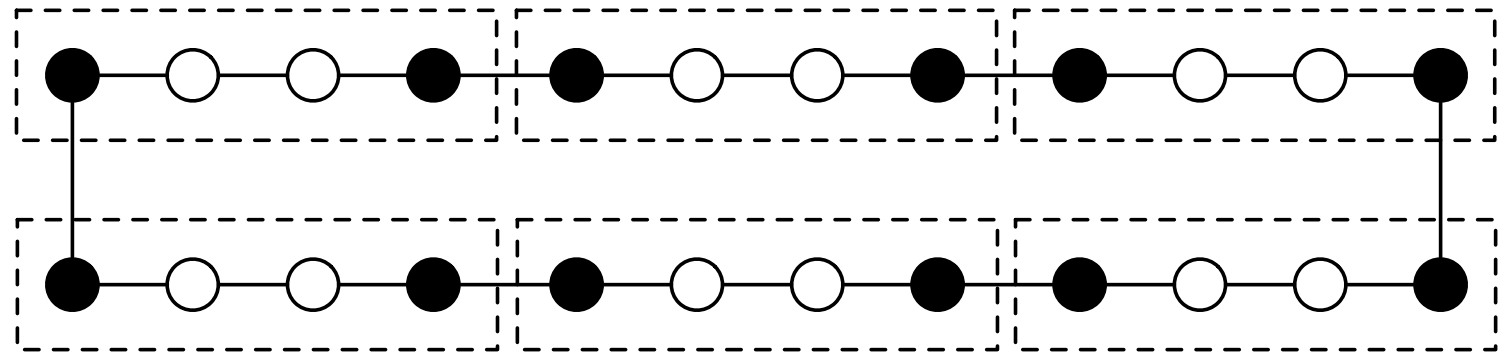
**For high controllability degree:**

1. Partition so that  $\|\Delta\|_{\infty}$  are small (weakly coupled subsystems)
2. Select local control nodes so that  $\|\Lambda\|_{\infty}$  is small (large local controllability)

With decoupled control, global controllability becomes a local property

# Examples

## Circulant graph



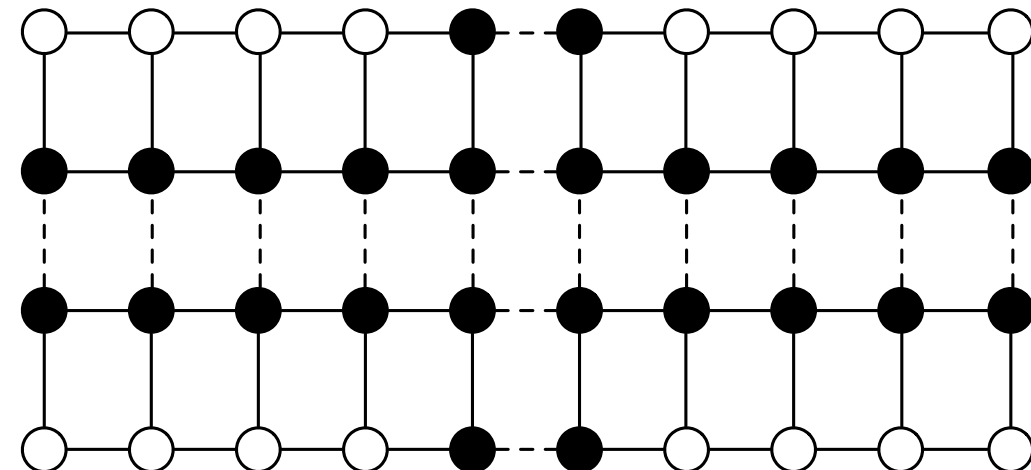
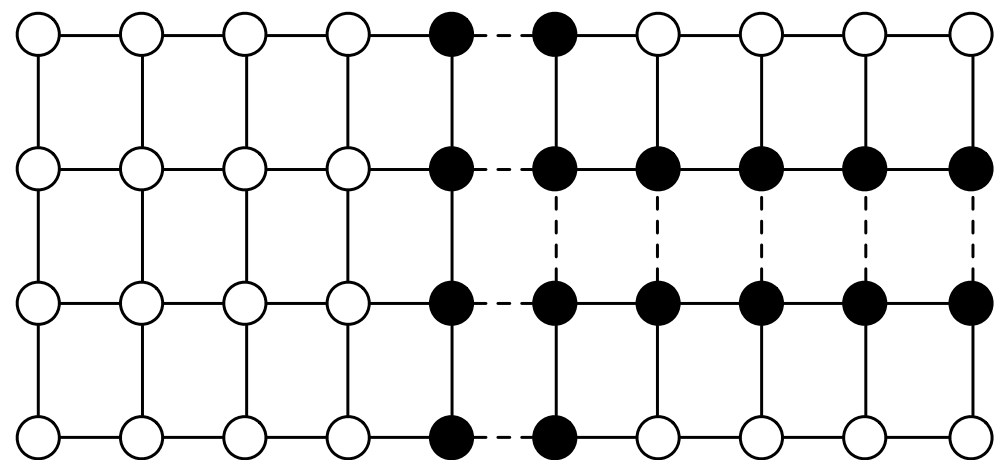
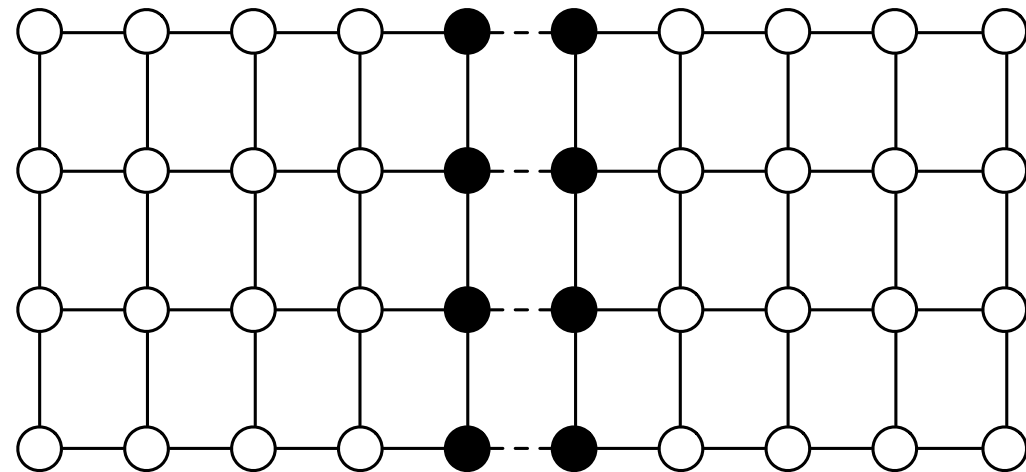
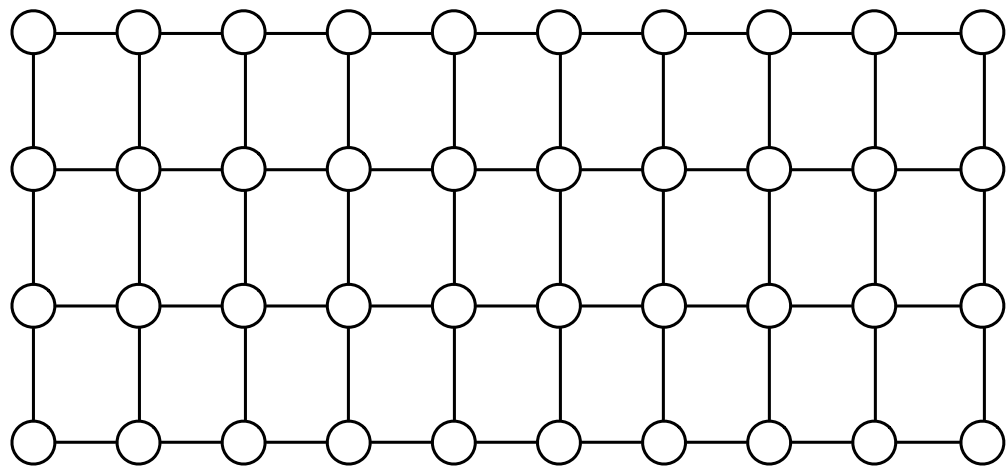
- $\lambda_{\min}$  with the decoupled control strategy
- - -  $\lambda_{\min}$  theoretical lower bound
- · -  $\lambda_{\min}$  with random positioning

# Network partition and selection of the control nodes

## Selection of the control nodes

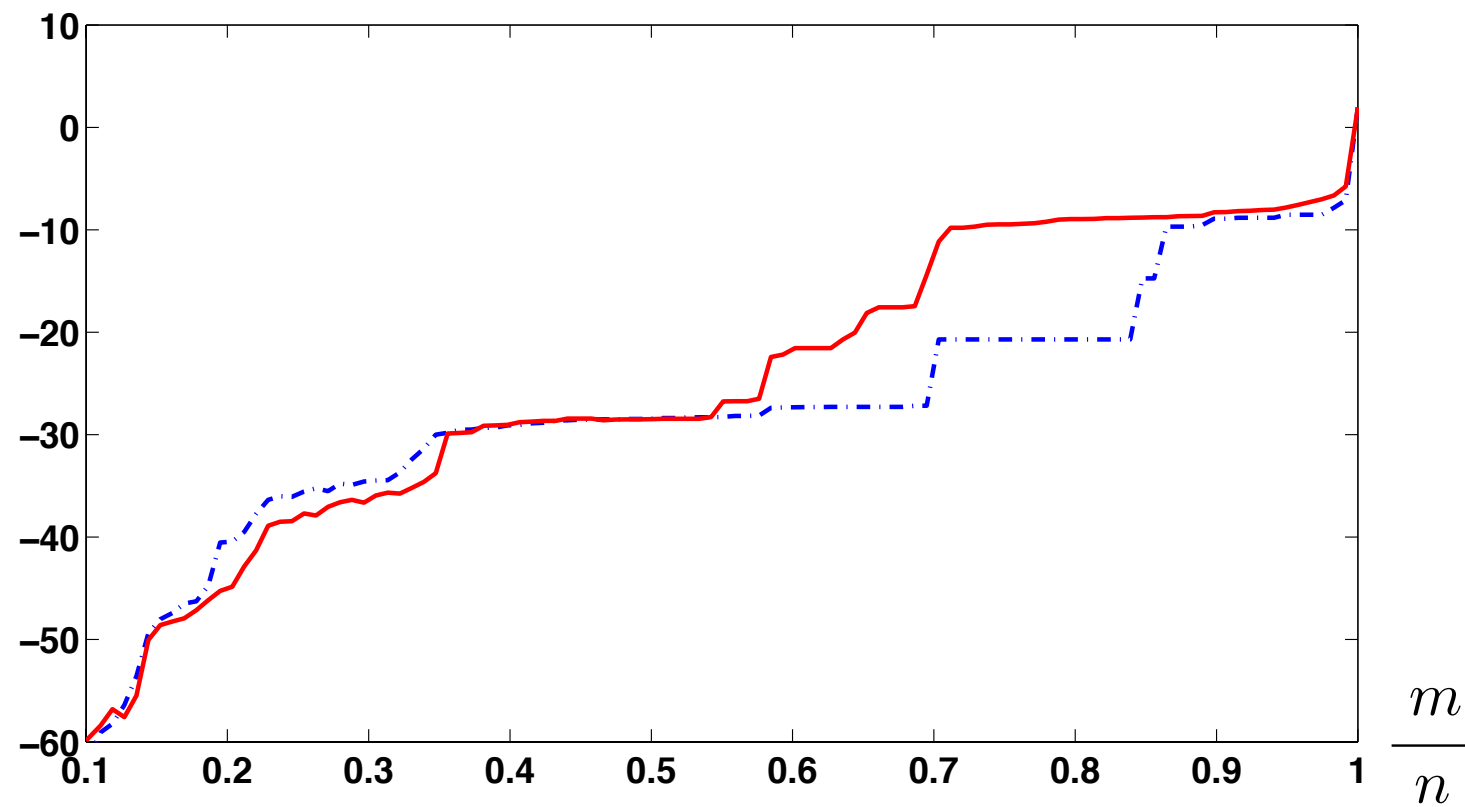
Until the desired number of control nodes have been selected:

1. Bisect the least controllable subsystem via Fiedler partitioning.
2. Include boundary nodes in the control set.



# Examples

## Power grid with 118 nodes



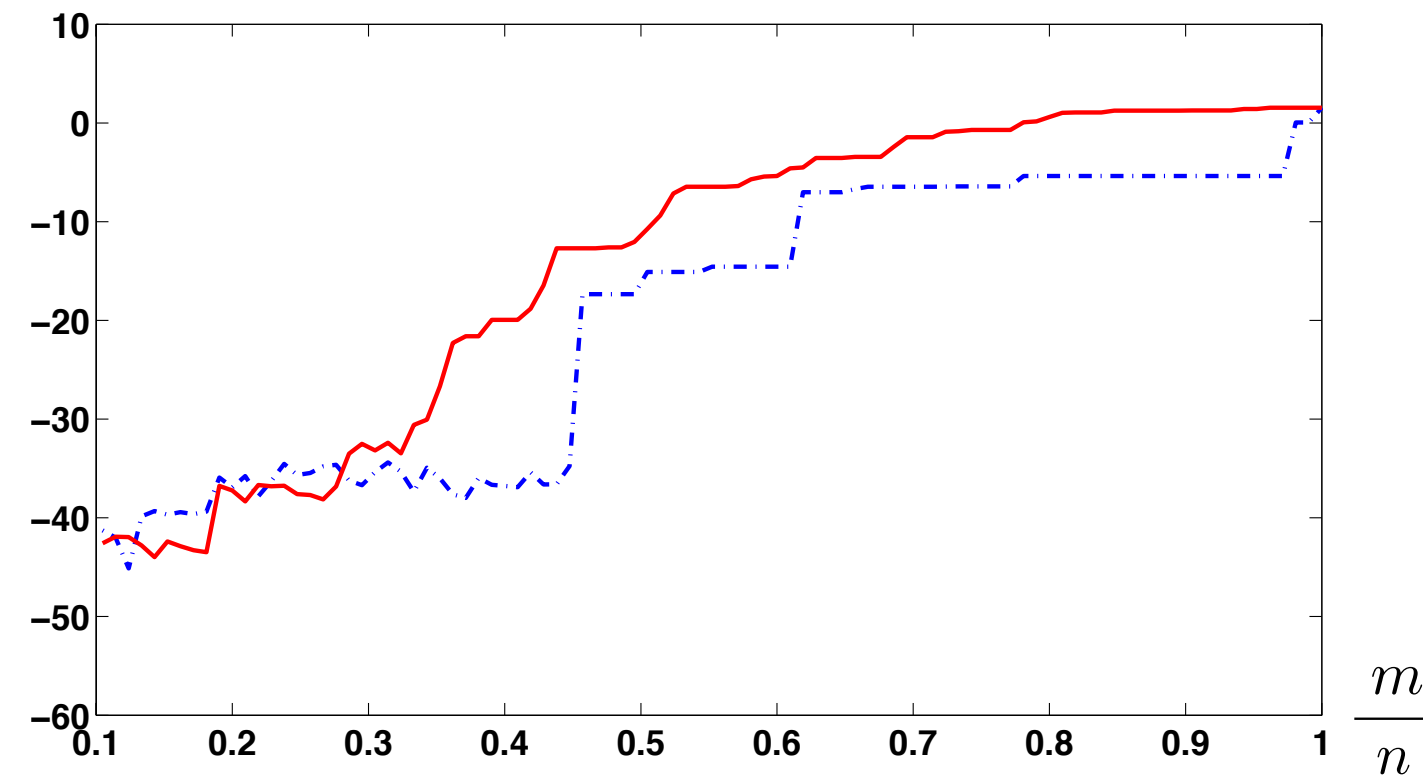
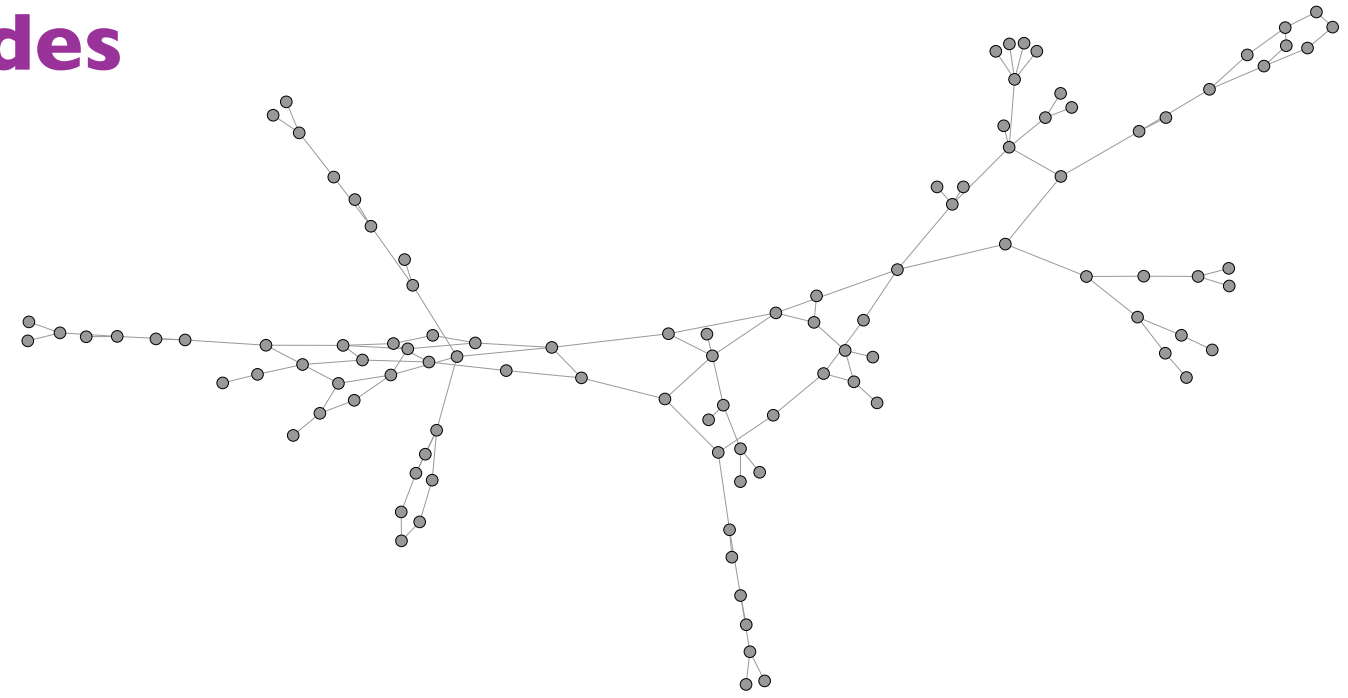
—  $\lambda_{\min}(W_T)$  with the decoupled control strategy

- · -  $\lambda_{\min}(W_T)$  with random positioning



# Examples

## Epidemics network with 86 nodes

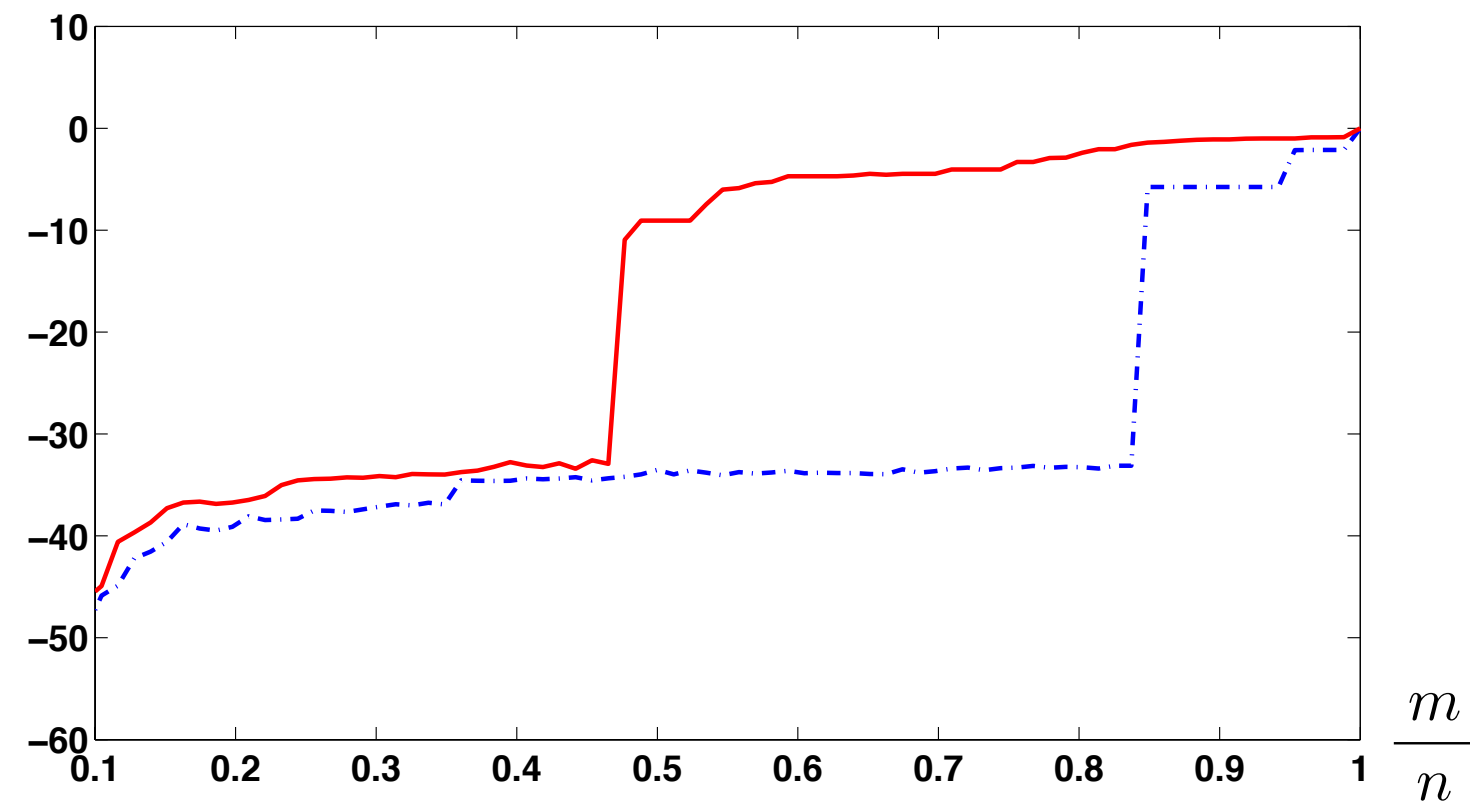
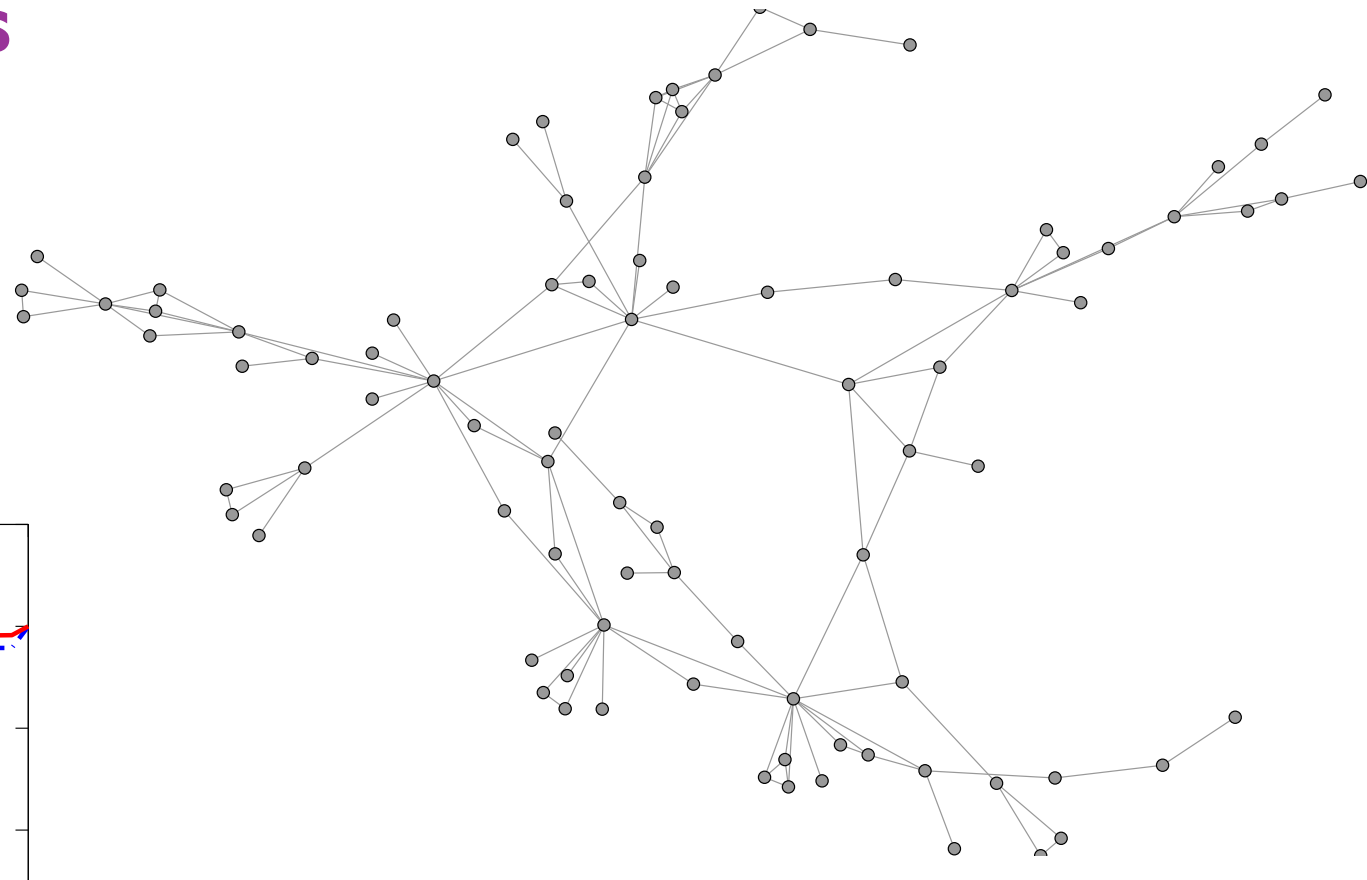


—  $\lambda_{\min}(W_T)$  with the decoupled control strategy

- - -  $\lambda_{\min}(W_T)$  with random positioning

# Examples

## Social network with 105 nodes



- $\lambda_{\min}(W_T)$  with the decoupled control strategy
- - -  $\lambda_{\min}(W_T)$  with random positioning

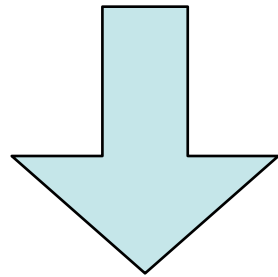


# Conclusions

- Similar results for **observability**
- For symmetric (isotropic) networks we need to control **a fixed fraction** of nodes
- For anisotropic networks it is enough to control **a fixed number** of nodes
- **Random** positioning works pretty well
- **Phase transition** can be noticed (critical fraction of controlled nodes)
- There are a lot of open problems:
  - Understanding isotropic and anisotropic networks
  - Controllability of random and of structured graphs
  - Performance of random positioning
  - Phase transition
  - Different metrics

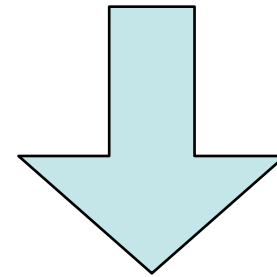
# Conclusions

**Controllability**



**Graph theory**

**Controllability degree**



**Spectral graph theory**



**Thank you**