

# Symbolic Control of Incrementally Stable Systems

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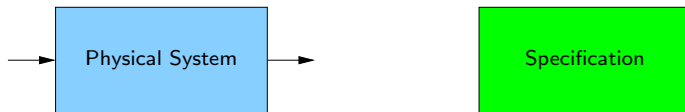
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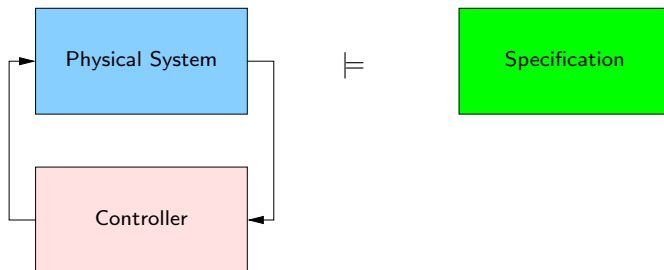
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LCCC, Lund, April 17-19 2013*



Algorithmic synthesis of controllers from high level specifications:



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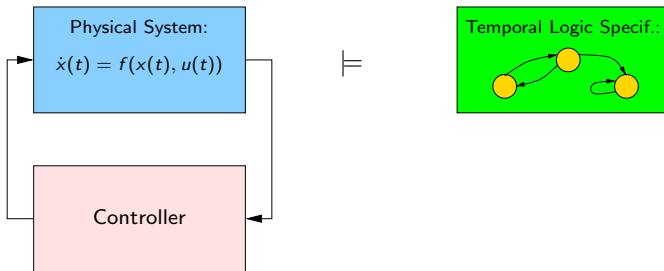
- Specifications can be expressed using temporal logic (e.g. LTL):

Safety	$\Box S$	(Always $S$ )
Reachability	$\Diamond T$	(Eventually $T$ )
Stability	$\Diamond(\Box T)$	
Recurrence	$\Box(\Diamond T)$	
Sequencing	$\Diamond(T_1 \wedge \Diamond T_2)$	
Coverage	$\Diamond T_1 \wedge \Diamond T_2$	
Fault recovery	$\Box(F \implies \Diamond R)$	

- LTL formula admits an equivalent (Büchi) automaton.

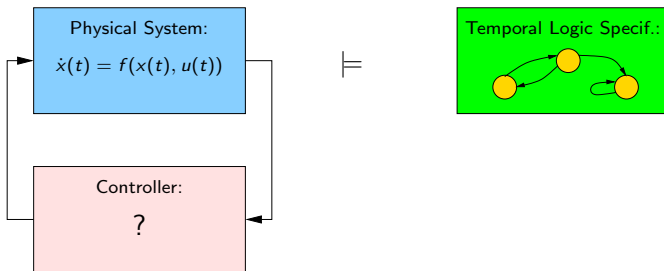
# Motivation

Algorithmic synthesis of controllers from high level specifications:



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Algorithmic synthesis of controllers from high level specifications:



The problem is hard because the model and the specification are heterogeneous.

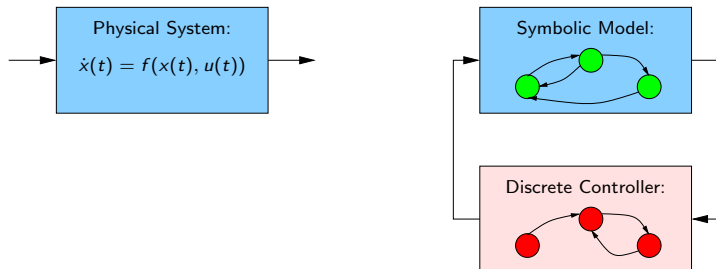
# Symbolic Approach to Control Synthesis

Approximate symbolic (*discrete*) model that is “formally related” to the (*continuous*) dynamics of the physical system:



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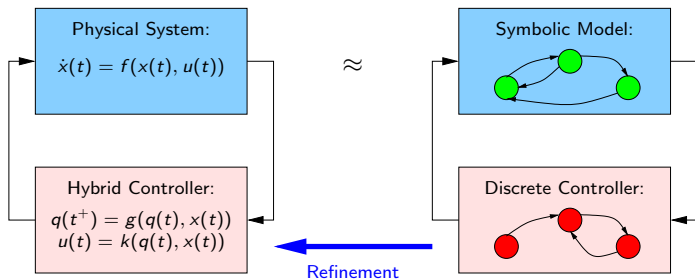
Approximate symbolic (*discrete*) model that is “formally related” to the (*continuous*) dynamics of the physical system:





# Symbolic Approach to Control Synthesis

Approximate symbolic (*discrete*) model that is “formally related” to the (*continuous*) dynamics of the physical system:



- 1 Behavioral metrics for discrete and continuous systems
  - Language metric
  - Approximate bisimulation and bisimulation metric
- 2 Symbolic abstractions of incrementally stable systems
  - Incrementally stable switched systems
  - State-space approaches: from uniform to multi-scale abstractions
  - Input-space approach

Unified modeling framework of discrete and (sampled) continuous systems.

## Definition

A transition system is a tuple  $T = (X, U, \delta, Y, H, X^0)$  where

- $X$  is a (discrete or continuous) set of states;
- $U$  is a (discrete or continuous) set of inputs;
- $\delta : X \times U \rightarrow 2^X$  is a transition relation;
- $Y$  is a (discrete or continuous) set of outputs;
- $H : X \rightarrow Y$  is an output map;
- $X^0 \subseteq X$  is a set of initial states.

The transition system is said to be *discrete* or *symbolic* if  $X$  and  $U$  are countable or finite.

- A *trajectory* of the transition system  $T$  is a finite or infinite sequence:

$$s = (x_0, u_0), (x_1, u_1), (x_2, u_2) \dots$$

where  $x_0 \in X^0$  and  $x_{k+1} \in \delta(x_k, u_k)$ ,  $\forall k$ .

- The associated *observed trajectory* is

$$o = (y_0, u_0), (y_1, u_1), (y_2, u_2) \dots \text{ where } y_k = H(x_k), \forall k.$$

- The set  $L(T)$  of observed trajectories of  $T$  is the *language* of transition system  $T$ .

- Traditional behavioral relationships for transition systems are based on language inclusion or equivalence.
- For systems observed over metric spaces, the distance between observed trajectories is more natural.
- Let  $T_i = (X_i, U, \delta_i, Y, H_i, X_i^0)$ ,  $i \in \{1, 2\}$ , be transition systems with a common set of inputs  $U$  and outputs  $Y$  equipped with a metric  $d$ . For  $o^1 \in L(T_1)$ ,  $o^2 \in L(T_2)$ ,

$$d(o^1, o^2) = \begin{cases} \sup_k d(y_k^1, y_k^2) & \text{if } u_k^1 = u_k^2, \forall k \\ +\infty & \text{otherwise} \end{cases}$$

## Definition

The language metric between  $T_1$  and  $T_2$  is given by

$$d_L(T_1, T_2) = \max \left\{ \sup_{o^1 \in L(T_1)} \inf_{o^2 \in L(T_2)} d(o^1, o^2), \sup_{o^2 \in L(T_2)} \inf_{o^1 \in L(T_1)} d(o^1, o^2) \right\}$$

- The language metric is generally hard to compute:
  - The choice of trajectory  $o^2$  approximating  $o^1$  may require knowledge of the whole trajectory  $o^1$ .
- Easier if the approximating trajectory can be selected transition after transition:
  - Bisimulation equivalence in the traditional setting.
  - Natural extension given by the bisimulation metric.

## Definition

Let  $\varepsilon \in \mathbb{R}_0^+$ , a relation  $R \subseteq X_1 \times X_2$  is an  $\varepsilon$ -*approximate bisimulation relation* if for all  $(x_1, x_2) \in R$  :

- 1  $d(H_1(x_1), H_2(x_2)) \leq \varepsilon$ ;
- 2  $\forall u \in U, \forall x'_1 \in \delta_1(x_1, u), \exists x'_2 \in \delta_2(x_2, u),$  such that  $(x'_1, x'_2) \in R$ ;
- 3  $\forall u \in U, \forall x'_2 \in \delta_2(x_2, u), \exists x'_1 \in \delta_1(x_1, u),$  such that  $(x'_1, x'_2) \in R$ .

## Definition

$T_1$  and  $T_2$  are  $\varepsilon$ -*approximately bisimilar* ( $T_1 \sim_\varepsilon T_2$ ) if :

- 1 For all  $x_1 \in X_1^0$ , there exists  $x_2 \in X_2^0$ , such that  $(x_1, x_2) \in R$ ;
- 2 For all  $x_2 \in X_2^0$ , there exists  $x_1 \in X_1^0$ , such that  $(x_1, x_2) \in R$ .

## Definition

The bisimulation metric between  $T_1$  and  $T_2$  is given by

$$d_B(T_1, T_2) = \inf \{ \varepsilon \in \mathbb{R}_0^+ \mid T_1 \sim_\varepsilon T_2 \}$$

- Fixed-point computation of the bisimulation metric for symbolic systems.
- For other systems, computation of upper-bounds using the notion of *bisimulation functions*.

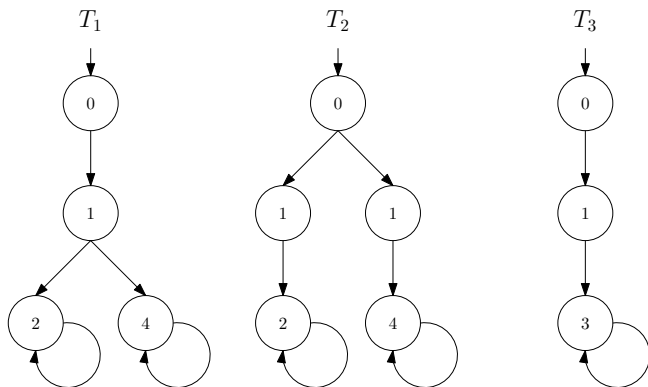
## Theorem

*The following inequality holds*

$$d_L(T_1, T_2) \leq d_B(T_1, T_2).$$



# A Simple Example



$$d_L(T_1, T_2) = 0, \quad d_B(T_1, T_2) = 2.$$

$$d_L(T_1, T_3) = 1, \quad d_B(T_1, T_3) = 1.$$

- 1 Behavioral metrics for discrete and continuous systems
  - Language metric
  - Approximate bisimulation and bisimulation metric
- 2 Symbolic abstractions of incrementally stable systems
  - Incrementally stable switched systems
  - State-space approaches: from uniform to multi-scale abstractions
  - Input-space approach

Continuous control systems with finite set of inputs:

## Definition

A switched system is a tuple  $\Sigma = (\mathbb{R}^n, P, \mathcal{F})$  where:

- $\mathbb{R}^n$  is the state space;
- $P = \{1, \dots, m\}$  is the finite set of modes;
- $F = \{f_p : \mathbb{R}^n \rightarrow \mathbb{R}^n \mid p \in P\}$  is the collection of vector fields.

For a switching signal  $\mathbf{p} : \mathbb{R}^+ \rightarrow P$ , initial state  $x \in \mathbb{R}^n$ ,  $\mathbf{x}(t, x, \mathbf{p})$  denotes the trajectory of  $\Sigma$  given by:

$$\dot{\mathbf{x}}(t) = f_{\mathbf{p}(t)}(\mathbf{x}(t)), \quad \mathbf{x}(0) = x.$$

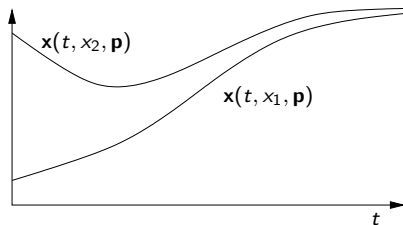
# Incremental Stability

Asymptotic forgetfulness of past history:

## Definition

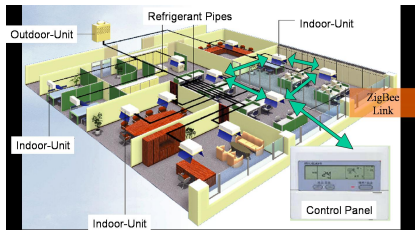
The switched system  $\Sigma$  is *incrementally globally uniformly asymptotically stable* ( $\delta$ -GUAS) if there exists a  $\mathcal{KL}$  function  $\beta$  such that for all initial conditions  $x_1, x_2 \in \mathbb{R}^n$ , for all switching signals  $\mathbf{p} : \mathbb{R}^+ \rightarrow P$ , for all  $t \in \mathbb{R}^+$ :

$$\|\mathbf{x}(t, x_1, \mathbf{p}) - \mathbf{x}(t, x_2, \mathbf{p})\| \leq \beta(\|x_1 - x_2\|, t) \rightarrow_{t \rightarrow +\infty} 0.$$



# Examples of incrementally stable systems

- Power converters.
- Thermal dynamics in buildings.
- Road traffic.



# Lyapunov Characterization

## Definition

$V : \mathbb{R}^n \times \mathbb{R}^n \rightarrow \mathbb{R}^+$  is a *common  $\delta$ -GUAS Lyapunov function* for  $\Sigma$  if there exist  $\mathcal{K}_\infty$  functions  $\underline{\alpha}$ ,  $\bar{\alpha}$  and  $\kappa \in \mathbb{R}^+$  such that for all  $x_1, x_2 \in \mathbb{R}^n$ :

$$\underline{\alpha}(\|x_1 - x_2\|) \leq V(x_1, x_2) \leq \bar{\alpha}(\|x_1 - x_2\|),$$

$$\forall p \in P, \quad \frac{\partial V}{\partial x_1}(x_1, x_2)f_p(x_1) + \frac{\partial V}{\partial x_2}(x_1, x_2)f_p(x_2) \leq -\kappa V(x_1, x_2).$$

## Theorem

*If there exists a common  $\delta$ -GUAS Lyapunov function, then  $\Sigma$  is  $\delta$ -GUAS.*

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## Theorem

*If there exists a common  $\delta$ -GUAS Lyapunov function, then  $\Sigma$  is  $\delta$ -GUAS.*

Supplementary assumption (true if working on a compact subset of  $\mathbb{R}^n$ ):

There exists a  $\mathcal{K}_\infty$  function  $\gamma$  such that

$$\forall x_1, x_2, x_3 \in \mathbb{R}^n, \quad |V(x_1, x_2) - V(x_1, x_3)| \leq \gamma(\|x_2 - x_3\|).$$

# Switched Systems as Transition Systems

- Consider a switched system  $\Sigma = (\mathbb{R}^n, P, \mathcal{F})$  and a time sampling parameter  $\tau > 0$ .
- Let  $T_\tau(\Sigma)$  be the transition system where:
  - the set of states is  $X = \mathbb{R}^n$ ;
  - the set of inputs is  $U = P$ ;
  - the transition relation is given by

$$x' \in \delta(x, p) \iff x' = \mathbf{x}(\tau, x, p);$$

- the set of outputs is  $Y = \mathbb{R}^n$ ;
- the output map  $H$  is the identity map over  $\mathbb{R}^n$ ;
- the set of initial states is  $X^0 = \mathbb{R}^n$ .



# Computation of the Symbolic Abstraction

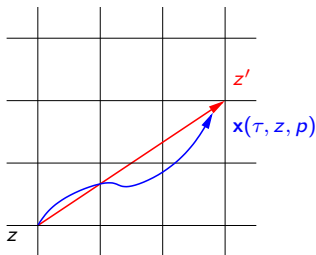
- We start by approximating the set of states  $\mathbb{R}^n$  by:

$$[\mathbb{R}^n]_\eta = \left\{ z \in \mathbb{R}^n \mid z_i = k_i \frac{2\eta}{\sqrt{n}}, k_i \in \mathbb{Z}, i = 1, \dots, n \right\},$$

where  $\eta > 0$  is a state sampling parameter:

$$\forall x \in \mathbb{R}^n, \exists z \in [\mathbb{R}^n]_\eta, \|x - z\| \leq \eta.$$

- Approximation of the transition relation = “rounding”:



# Approximation Theorem

## Theorem

Let us assume that there exists  $V : \mathbb{R}^n \times \mathbb{R}^n \rightarrow \mathbb{R}^+$  which is a common  $\delta$ -GUAS Lyapunov function for  $\Sigma$ . Consider sampling parameters  $\tau, \eta \in \mathbb{R}^+$  and a desired precision  $\varepsilon \in \mathbb{R}^+$ . If

$$\eta \leq \min \left\{ \gamma^{-1} \left( (1 - e^{-\kappa\tau}) \underline{\alpha}(\varepsilon) \right), \bar{\alpha}^{-1}(\underline{\alpha}(\varepsilon)) \right\}$$

then, the relation  $R \subseteq \mathbb{R}^n \times [\mathbb{R}^n]_\eta$  given by

$$R = \left\{ (x, z) \in \mathbb{R}^n \times [\mathbb{R}^n]_\eta \mid V(x, z) \leq \underline{\alpha}(\varepsilon) \right\}$$

is an  $\varepsilon$ -approximate bisimulation relation and  $T_\tau(\Sigma) \sim_\varepsilon T_{\tau, \eta}(\Sigma)$ .

Main idea of the proof: show that accumulation of successive “rounding errors” is contained by incremental stability.

# Comments on the Approximation Theorem

- Based on sampling (gridding) of time and space: simple to compute.
- For a given time sampling parameter  $\tau$ , any precision  $\varepsilon$  can be achieved by choosing appropriately the state sampling parameter  $\eta$  (the smaller  $\tau$  or  $\varepsilon$ , the smaller  $\eta$ ).
- Uniform time and space discretization: excessive computation time and memory consumption.
- Overcome this problem with multi-scale symbolic abstractions: on-the-fly refinement where fast switching needed, guided by controller synthesis.

# Switched Systems in a Multi-Scale Setting

- Consider a switched system  $\Sigma = (\mathbb{R}^n, P, \mathcal{F})$ , time and scale sampling parameters  $\tau > 0$  and  $N \in \mathbb{N}$ .
- We change the control paradigm: the (aperiodic) controller chooses a mode and a duration during which it will be applied.
- Let  $T_\tau^N(\Sigma)$  be the transition system where:
  - the set of states is  $X = \mathbb{R}^n$ ;
  - the set of inputs is  $U = P \times \Theta_\tau^N$  where  $\Theta_\tau^N = \{2^{-s}\tau \mid s = 0, \dots, N\}$ ;
  - the transition relation is given by

$$x' \in \delta(x, (p, 2^{-s}\tau)) \iff x' = \mathbf{x}(2^{-s}\tau, x, p);$$

- the set of outputs is  $Y = \mathbb{R}^n$ ;
- the output map  $H$  is the identity map over  $\mathbb{R}^n$ ;
- the set of initial states is  $X^0 = \mathbb{R}^n$ .

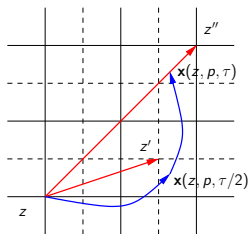
# Multi-Scale Symbolic Abstraction

- The set of states  $\mathbb{R}^n$  is approximated by a sequence of embedded lattices  $Q^0 \subseteq Q^1 \subseteq \dots \subseteq Q^N \subseteq \mathbb{R}^n$  with:

$$Q^s = [\mathbb{R}^n]_{2^{-s}\eta} = \left\{ z \in \mathbb{R}^n \mid z_i = k_i \frac{2^{-s+1}\eta}{\sqrt{n}}, k_i \in \mathbb{Z}, i = 1, \dots, n \right\}$$

where  $\eta > 0$  is a state sampling parameter:

- Approximation of the transition relation:



Fine scales reached only by transitions of shorter duration.

## Theorem

Let us assume that there exists  $V : \mathbb{R}^n \times \mathbb{R}^n \rightarrow \mathbb{R}^+$  which is a common  $\delta$ -GUAS Lyapunov function for  $\Sigma$ . Consider sampling and scale parameters  $\tau, \eta \in \mathbb{R}^+$ ,  $N \in \mathbb{N}$  and a desired precision  $\varepsilon \in \mathbb{R}^+$ . If

$$\eta \leq \min \left\{ \min_{s=0 \dots N} \left[ 2^s \gamma^{-1} \left( (1 - e^{-\kappa 2^{-s} \tau}) \underline{\alpha}(\varepsilon) \right) \right], \bar{\alpha}^{-1}(\underline{\alpha}(\varepsilon)) \right\}$$

then, the relation  $R \subseteq \mathbb{R}^n \times Q^N$  given by

$$R = \left\{ (x, z) \in \mathbb{R}^n \times Q^N \mid V(x, z) \leq \underline{\alpha}(\varepsilon) \right\}$$

is an  $\varepsilon$ -approximate bisimulation relation and  $T_\tau(\Sigma) \sim_\varepsilon T_{\tau, \eta}(\Sigma)$ .

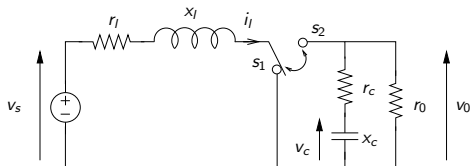
# Controller Synthesis using Multi-Scale Abstractions

- Multi-scale abstractions are computed on the fly during controller synthesis using depth first search algorithm:
  - Start from initial states:
    - elements of the coarsest lattice.
  - Explore transitions of longer duration first and transitions of shorter duration only if specification cannot be met by transitions of longer durations:
    - fine lattices are explored only when necessary.
- For safety specifications: notion of maximal lazy safety controller.
- Tool CoSyMA: Controller Synthesis using Multi-Scale abstractions.  
`multiscale-dcs.gforge.inria.fr`

# Example: DC-DC Converter

Power converter with switching control:

- Incrementally stable.
- Safety specification:  
[1.15, 1.55]  $\times$  [5.45, 5.85].



	Uniform abstraction $T_{\tau, \eta}(\Sigma)$ $\tau = 0.5, \eta = 0.0003, \varepsilon = 0.05$
Time	9.2s
Size ( $10^3$ )	936
Cont. ratio	93%

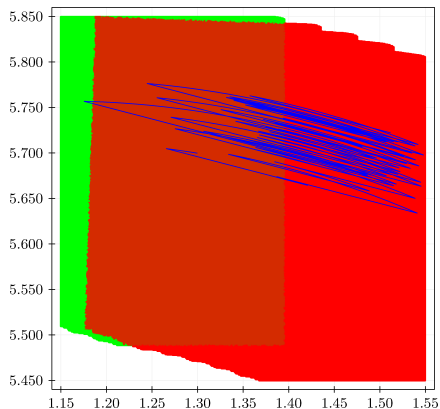
  

	Multi-scale abstraction $T_{\tau, \eta}^N(\Sigma)$ $N = 6, \tau = 32, \eta = 0.018, \varepsilon = 0.05$
Time	0.6s
Size ( $10^3$ )	6
Durations	4 (33%), 2 (9%), 1 (50%), 0.5 (8%)
Cont. Ratio	92%



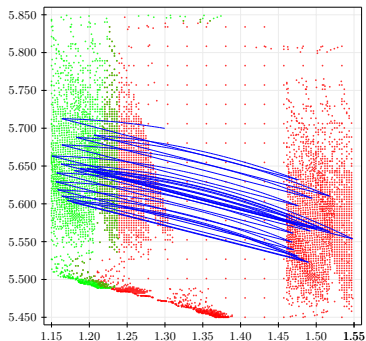
# Example: Boost DC-DC Converter

Uniform abstraction  $T_{\tau,\eta}(\Sigma)$ :

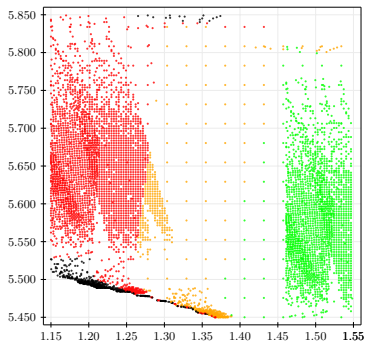


# Example: Boost DC-DC Converter

Multiscale abstraction  $\mathcal{T}_{\tau,\eta}^N(\Sigma)$ :



Modes

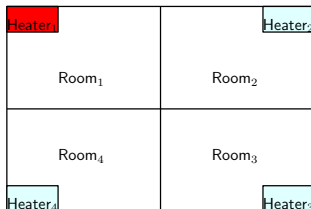


Durations

# Example: 4 Room Building

4 dimensional thermal model:

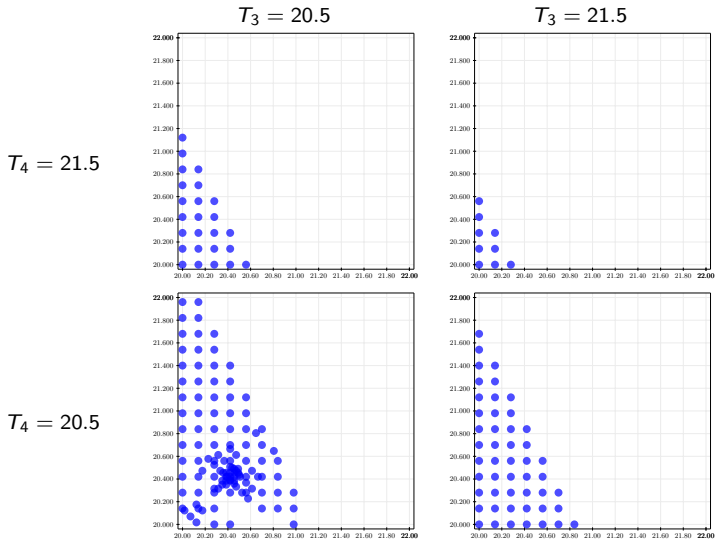
- Incrementally stable.
- At most one heater on at every instant.
- Safety specification:  $[20, 22]^4$ .



	Multi-scale abstractions $T_{\tau, \eta}^N(\Sigma)$ $N = 4, \tau = 80, \eta = 0.14, \varepsilon = 0.2$
Time	39s
Size ( $10^3$ )	232
Durations	20 (2%), 10 (91%), 5 (7%)
Cont. Ratio	99%

# Example: 4 Room Building

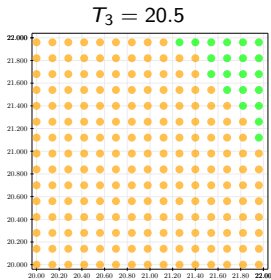
Control maps (mode 1):



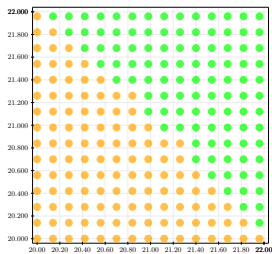
# Example: 4 Room Building

Control maps (durations):

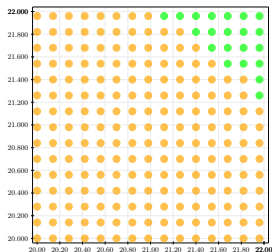
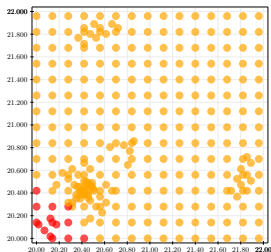
$T_4 = 21.5$



$T_3 = 21.5$



$T_4 = 20.5$



# Mode Sequences as Symbolic States

- State-space approaches suffer from the curse of dimensionality.
- Alternative: input-space approach
  - Incremental stability = asymptotic forgetfulness of past history,
  - Use mode sequences of given length  $N$ , representing the latest applied modes, as symbolic states of symbolic model  $T_{\tau, N}(\Sigma)$ ,
  - The transition relation is given for  $w = p_1 p_2 \dots p_n$  and  $p \in P$  by

$$w' \in \delta(w, p) \iff w' = p_2 \dots p_n p.$$

- The output map is defined for  $w = p_1 p_2 \dots p_n$  as

$$H(w) = \mathbf{x}(N\tau, x_s, \mathbf{p}_w) \text{ where } \mathbf{p}_w(t) = p_i, \forall t \in [(i-1)\tau, i\tau).$$

where  $x_s \in \mathbb{R}^n$  is a source state.

## Theorem

Let us assume that there exists  $V : \mathbb{R}^n \times \mathbb{R}^n \rightarrow \mathbb{R}^+$  which is a common  $\delta$ -GUAS Lyapunov function for  $\Sigma$ . Consider time sampling parameter  $\tau \in \mathbb{R}^+$ , sequence length  $N \in \mathbb{N}$  and a desired precision  $\varepsilon \in \mathbb{R}^+$ . Let

$$\varepsilon \geq \underline{\alpha}^{-1} \left( \frac{\gamma(e^{-N\kappa\tau}\theta(x_s))}{1 - e^{-\kappa\tau}} \right)$$

where  $\theta(x_s) = \max_{p \in P} V(\mathbf{x}(\tau, x_s, p), x_s)$ . Then, the relation  $R \subseteq \mathbb{R}^n \times P^N$  given by

$$R = \left\{ (x, w) \in \mathbb{R}^n \times P^N \mid V(x, H(w)) \leq \underline{\alpha}(\varepsilon) \right\}$$

is an  $\varepsilon$ -approximate bisimulation relation between  $T_\tau(\Sigma)$  and  $T_{\tau,\eta}(\Sigma)$ .

# Comments on the Approximation Theorem

- The source state can be chosen so as to minimize  $\theta(x_s)$ .
- For a given time sampling parameter  $\tau$ , any precision  $\varepsilon$  can be achieved by choosing appropriately the sequence length  $N$ .
- Number of symbolic states grows exponentially with the sequence length  $N$ .
- Asymptotic estimates show that for a given precision  $\varepsilon$ , the input-space approach leads to a smaller number of symbolic states than the (uniform) state-space approach as soon as

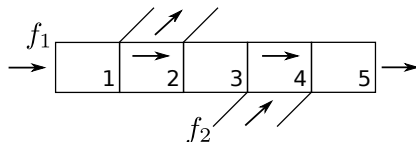
$$\ln(|P|) \leq \kappa T n.$$



# Example: Road Traffic

5 dimensional model:

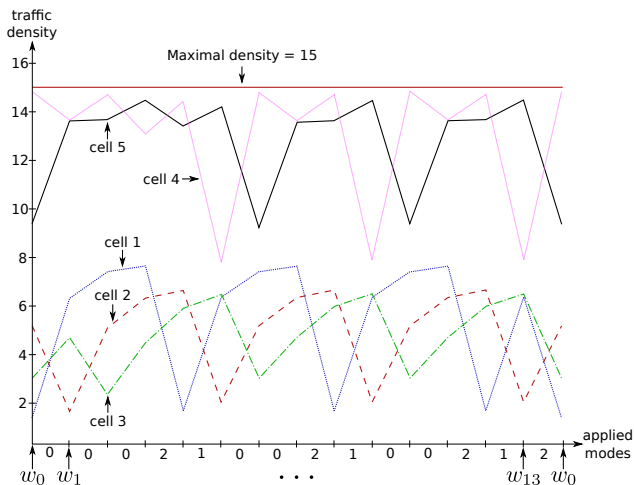
- Incrementally stable.
- At least one green light.
- Safety specification:  $[0, 15]^5$ .
- Fairness constraint: red light no longer than 3 time units.



Sequence length $N$	10	12	14
Size ( $10^3$ )	59	531	4783
Precision $\varepsilon$	0.1	0.01	0.001

# Example: Road Traffic

Periodic schedule for light coordination:



- Approximately bisimilar symbolic abstractions:
  - A rigorous tool for controller synthesis:  
Synthesized controllers are “correct by design” .
  - Allow us to leverage efficient algorithmic techniques from discrete systems to continuous and hybrid systems.
  - Computable for interesting classes of systems: switched systems, continuous control systems...
  - Several approaches can help to reduce the computation burden.
- Ongoing and future work:
  - Tool CoSyMA: Controller Synthesis using Multi-scale Abstractions.
  - Multi-scale input-space approaches.
  - Symbolic models for infinite dimensional systems.

- Behavioral metrics for discrete and continuous systems:
  - G. and Pappas, Approximation metrics for discrete and continuous systems. IEEE TAC, 2007.
- Symbolic abstractions of incrementally stable systems:
  - Pola, G. and Tabuada, Approximately bisimilar symbolic models for nonlinear control systems. Automatica, 2008.
  - G., Pola and Tabuada, Approximately bisimilar symbolic models for incrementally stable switched systems. IEEE TAC, 2010.
  - Camara, G. and Goessler, Safety controller synthesis for switched systems using multi-scale symbolic models. CDC, 2011.
  - Mouelhi, G. and Goessler, CoSyMA: a tool for controller synthesis using multi-scale abstractions. HSCC, 2013.
  - Le Corronc, G. and Goessler, Mode sequences as symbolic states in abstractions of incrementally stable switched systems. Submitted.