

multiform

Integrated Multi-formalism Tool Support for the Design of Networked Embedded Control Systems

Bridging between different modeling formalisms - results from the MULTIFORM project

Martin Hüfner, Christian Sonntag, Sebastian Engell

Process Dynamics and Operations Group
Department of Biochemical and Chemical Engineering
TU Dortmund
Germany

Outline

- The MULTIFORM project
- Design flow example
- Tool developments
- Model exchange and model transformations
- Lessons learned



MULTIFORM: EU ICT STREP 9/2008 - 5/2012



VEMAC

Aachen, Germany
 Michael Reke

KVCA

- "Danish Cooling Cluster"
 Jens Andersen
- Closely working with DANFOSS

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TUDO (Coordinator)

TU Dortmund, Germany
 Sebastian Engell

TUE

 TU Eindhoven,
 Netherlands
 Koos Rooda, Bert van Beek, Jos Baeten

Verimag/UJF

Universite Joseph

 Fourier, Grenoble, France
 Goran Frehse, Oded
 Maler

RWTH

RWTH Aachen, Germany
 Stefan Kowalewski

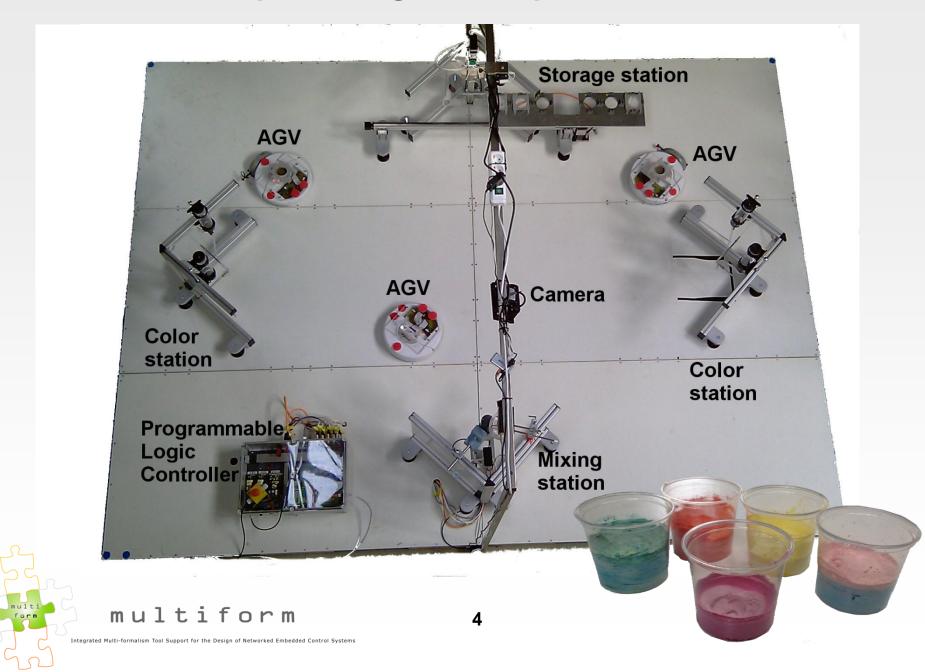
AAU

Aalborg Universitet,
 Denmark
 Kim Larsen, Brian Nielsen

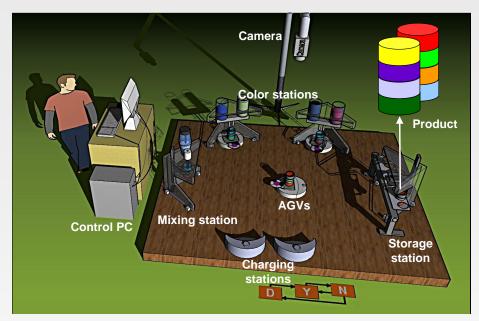
ESI

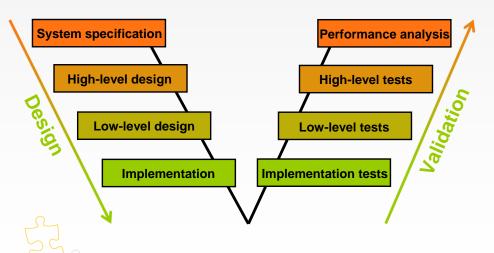
Stichting Embedded
 Systems Institute
 Ed Brinksma, Boudewijn
 Haverkort

Example: Design of a Pipeless Plant



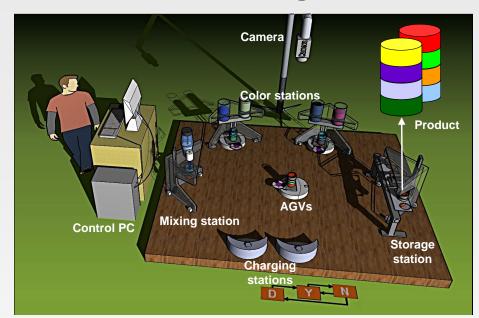
Challenges for Model-based Design (1)





- Design and validation on different levels of abstraction
 - Specification
 - Specification of the tasks and of the performance of the system
 - High-level design
 - Choice of the equipment, feasibility and bottleneck analysis, throughput maximization, plant layout optimization
 - Low-level design
 - Optimization and control of processing steps and motion dynamics, logic control
 - Choice of sensors and actuators, communication system
 - Implementation
 - PLCs, embedded controllers, communication system

Challenges for Model-based Design (2)



The control system spans the complete control hierarchy

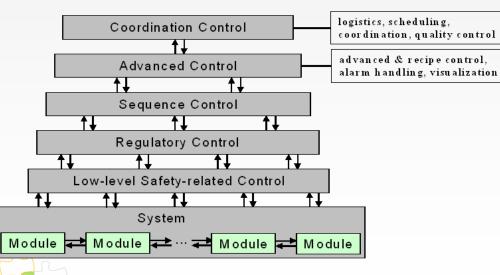
Coordination control Timed or hybrid models

- Scheduling and performance optimization
- Advanced control
 - Control of batch processes
 - AGV path planning Continuous models
- Regulatory control
 - AGV motion control
- hybrid, and continuous

Discrete-event,

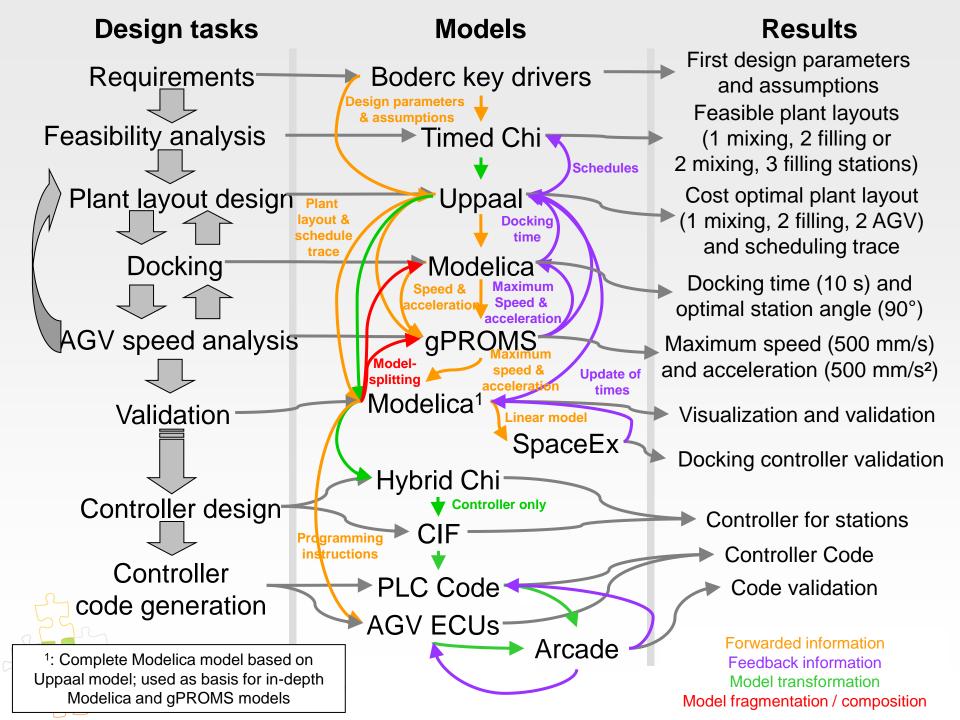
- Docking control
- Sequence control in the processing stations
- Low-level continuous control
- Low-level safety-related control

Discrete-event, timed, and hybrid models





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Integrated Model-based Design

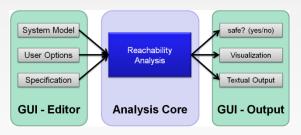
- Integrated modeling and design of the system itself and of the multilayered and networked control system
 - Including a structured approach to the management of specifications, design decisions, models, and results
- Coverage of all layers of the automation and design hierarchy
 - Integrated tool support on all layers of the automation and design hierarchies
 - Current state: Islands of support for specific design and analysis tasks
- Trans-level integration of model-based design approaches
- Support of iterations in the design process
 - Propagation of faults and unexpected behaviors
 - Modifications over the life cycle without top-down redesign
- Improvement of the tool support for the design steps
- Tool integration and Design Framework
- Exchange of models between tools via the CIF (Compositional Interchange Format)

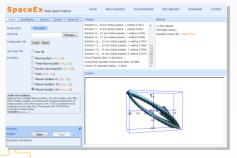


MULTIFORM Tools and Tool Chains

Verification

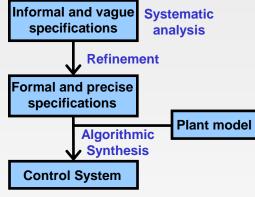
- Verification tool SpaceEx (successor of PHAVer)
- Consistency checking methods using UPPAAL
- Step-wise refinement based on HCIF



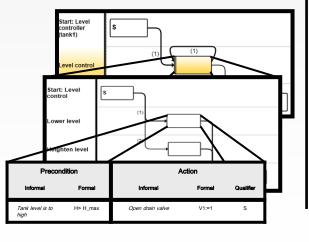


Logic Controller Design

Integrated controller design and analysis

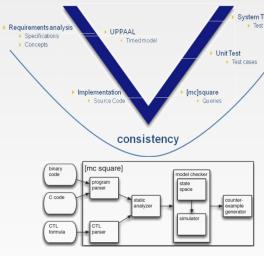


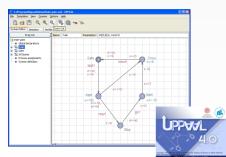
Specification using DC/FT



Code Analysis

Code and requirements analysis for ECUs using Arcade and UPPAAL

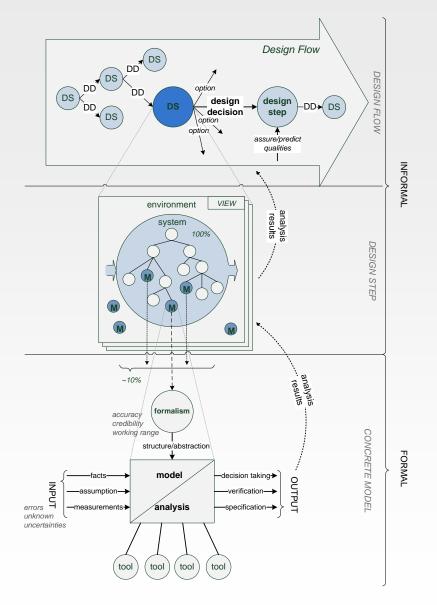






The MULTIFORM Design Framework [ESI]

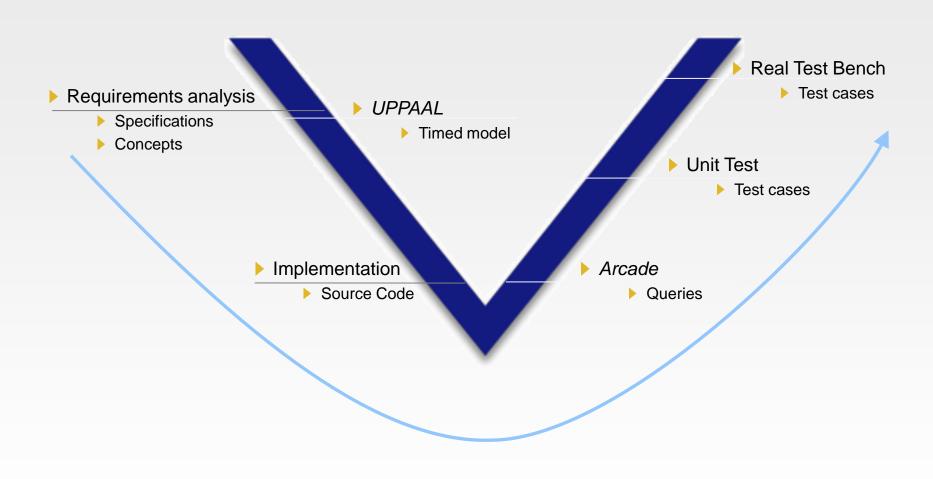
- Consistent integration of design models into a common software framework
- Support of a generic design flow model
 - Design decisions
 - System design
- Consistency management
 - Communication of design parameters
 - Conflict detection
 - Models and results management





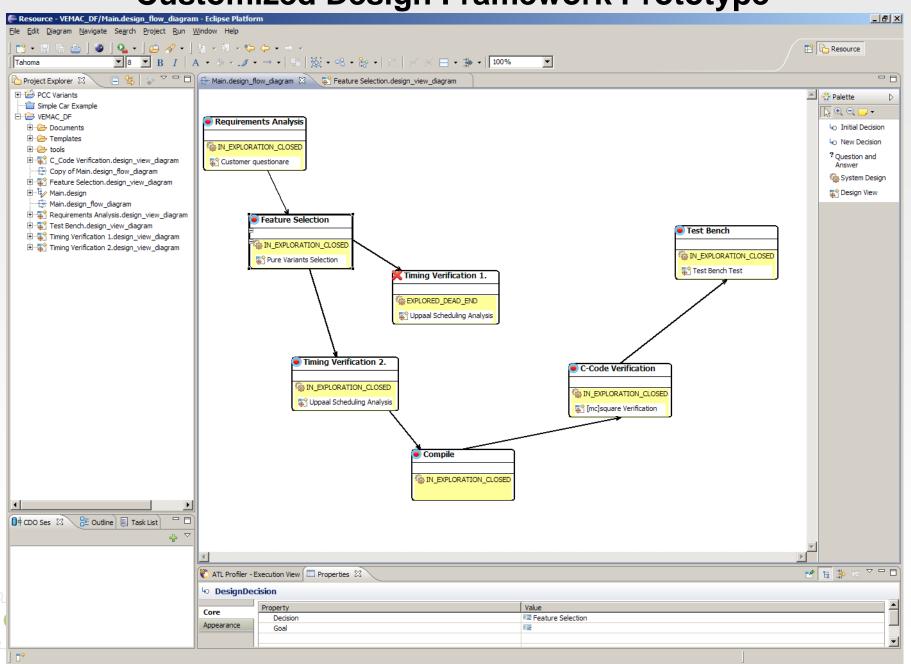
VEMAC Development Process

V-Model



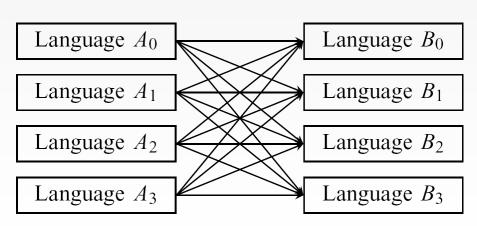


Customized Design Framework Prototype



Model Exchange Using the Compositional Interchange Format (CIF)

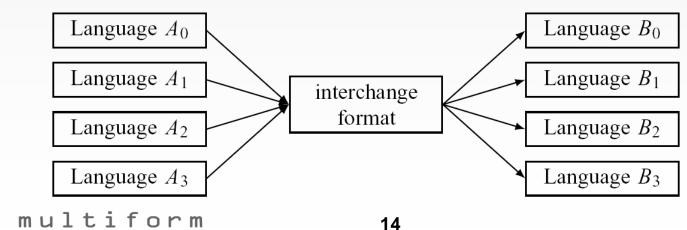
- Incompatibility of tools is one of the major obstacles for a broader acceptance of model-based design in industry
- → Achieve inter-operability by (algorithmic) model transformations
- One possibility: Bi-lateral transformations
 - Problems
 - Many transformations may be needed
 - The developer of a transformation must be familiar with many different formalisms





Model Exchange with the Compositional Interchange Format (CIF)

- Incompatibility of tools is one of the major obstacles for a broader acceptance of model-based design in industry
- → Achieve inter-operability by (algorithmic) model transformations
- One possibility: Bi-lateral transformations
- Interchange format
 - Generic and sufficiently rich modelling formalism
 - Only transformations from/to the interchange format are necessary
 - → Reduction of the implementation effort



egrated Multi-formalism Tool Support for the Design of Networked Embedded Control Systems

The Compositional Interchange Format (CIF) [Bert van Beek et al., TUE]

- Purposes
 - Establish inter-operability of a wide range of tools
 - Provide a generic formalism for general hybrid systems
- Major features
 - Formal and compositional semantics
 - Independent of implementation aspects
 - Mathematical correctness proofs of translations
 - → Property-preserving model transformations possible
 - Fully implicit DAE dynamics (possibly discontinuous)
 - Hierarchy and re-usability
 - Parallel composition with different communication concepts
 - Model component interaction
 - Point to point communication, multi-component synchronization, broadcast communication, shared variables
 - Different urgency concepts





The Compositional Interchange Format (CIF) [Bert van Beek et al., TUE]

Transition

State

Invariants
(equations that are active when state is active)
e.g.:

v' = -q

Guards

(transition can only be taken if guard is true)

e.g.: **a > b**

Updates

(new discrete values or reinitialization)

e.g.: z := 5, $\{v\}$: new(v) = 2Synchronization

(between different automata via labels or channels)

Urgency

(nondeterminism, determinism, stochastic)

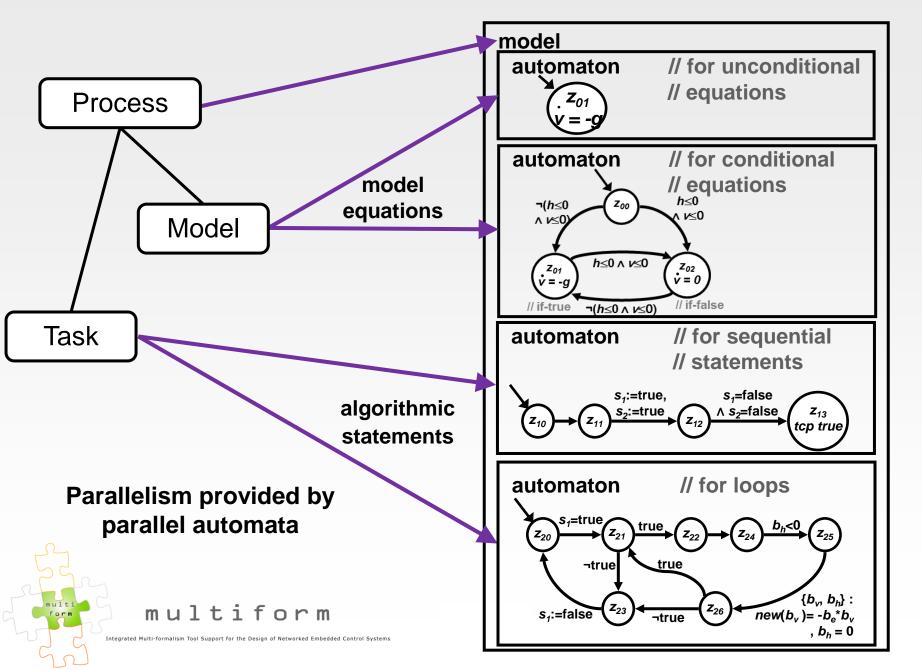
Initial
(Conditions if state is initially active)

Formal definition by Structural Operational Semantics (SOS) rules, e.g.:

$$\frac{(\alpha_0, \sigma) \xrightarrow{a, \text{true}, X} (\alpha'_0, \sigma'), (\alpha_1, \sigma) \xrightarrow{a, \text{true}, X} (\alpha'_1, \sigma')}{(\alpha_0 \parallel \alpha_1, \sigma) \xrightarrow{a, \text{true}, X} (\alpha'_0 \parallel \alpha'_1, \sigma')}$$



Transformations – gPROMS → CIF (Excerpt)



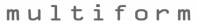
Simple Example: tank.cif

```
model TankController()=
| cont control real V = 10.0
           real Qi, Qo
: var
                                                           V
; disc control nat n = 0
:: Tank : |( mode physics = initial
             inv V' = Qi - Qo
              , Qi = n * 5.0
              , Qo = sqrt(V)
 Controller : |( mode closed = initial
                     (when V \le 2 now do n := 1) goto opened
             , opened = (when \lor >= 10 now do n := 0) goto closed
          )|
```



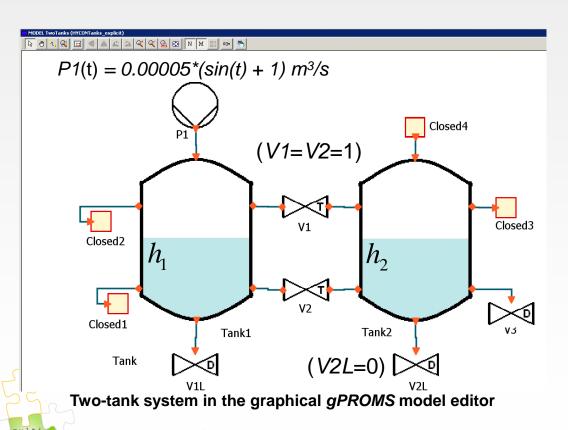
Flattened Example: tank_flat.cif

```
model TankController() =
| | var real V = 10.0; var real Qi; var real Qo; var nat n = 0
:: Tank Controller:
   var string Controller_LP ; var string Tank_LP
 ; controlset Controller LP, Tank LP, V, n
 ; dyntypemap disc Controller_LP; disc Tank_LP; disc n; cont V;
  mode X =
   initial (((Tank_LP) = ("physics")) and (true))
             and (((Controller LP) = ("closed")) and (true))
   inv ((Tank LP) = ("physics")) => ((((V)') = ((Qi) - (Qo)))
             and (((Qi) = ((n) * (5.0))) and ((Qo) = (sqrt(V))))
   tcp ((Controller_LP) = ("closed")) => (not ((V) <= (2)))
   tcp ((Controller_LP) = ("opened")) \Rightarrow (not ((V) \Rightarrow (10)))
    ( when (V) \le (2), (Controller LP) = ("closed") do
       {Controller_LP, n}: (new(n)) = (1), (new(Controller_LP)) = ("opened") )
    ( when (V) >= (10), (Controller_LP) = ("opened") do
       {Controller\_LP, n} : (new(n)) = (0), (new(Controller\_LP)) = ("closed") ) goto
   Χ
```



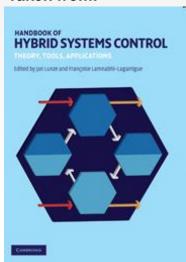
A Two-tank System under Discrete Control

- Hybrid non-linear model of a two-tank system, modeled in gPROMS
 - Designed to contain many constructs of the gPROMS language



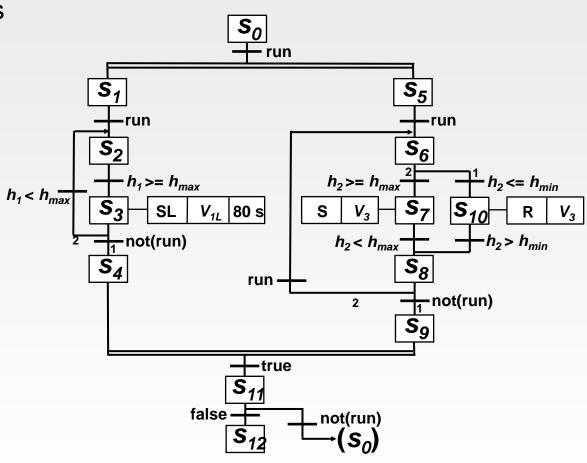
- Controlled variables:h₁, h₂
- Manipulated (discrete)variables: V1L, V3

Taken from:



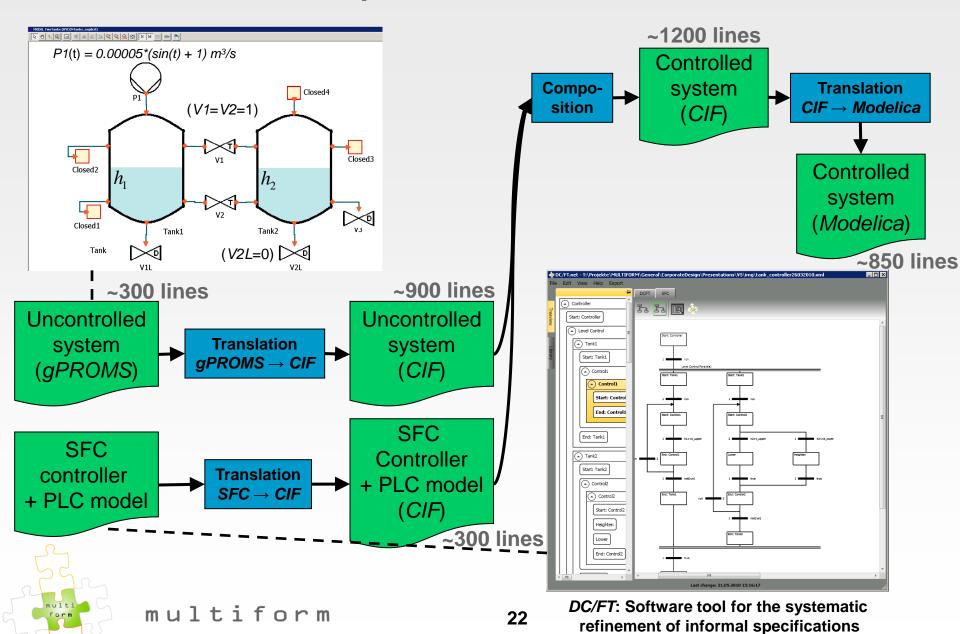
Two-tank Example: SFC Controller

- The SFC controller keeps the filling levels h_1 and h_2 between h_{min} =0.2 m and h_{max} =0.5 m
- If h₁ exceeds h_{max}, valve
 V1L is opened for 80s
- If h_2 exceeds h_{max} , valve V3 is opened until h_2 falls below h_{min}





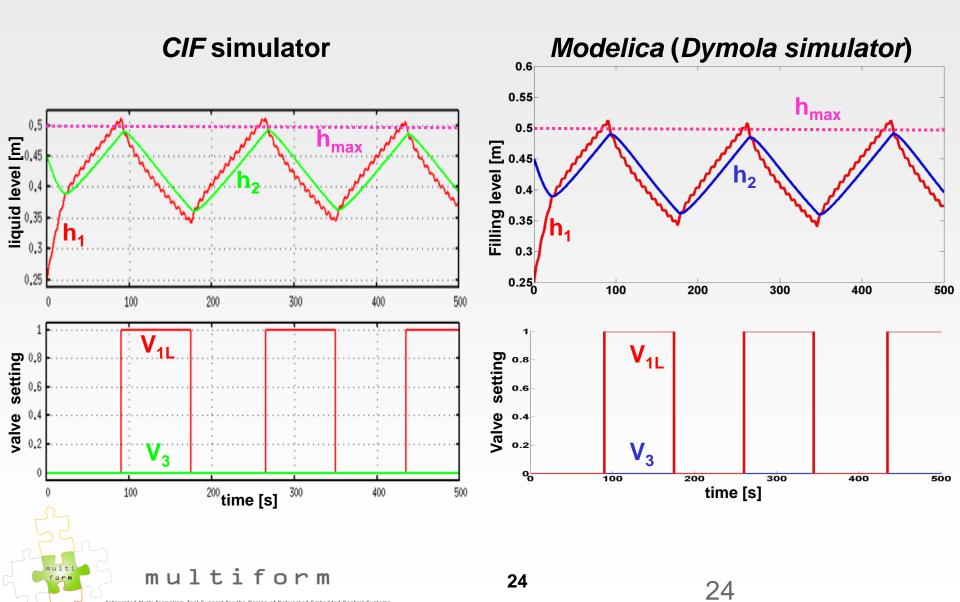
Two-tank Example: Transformation Tool Chain



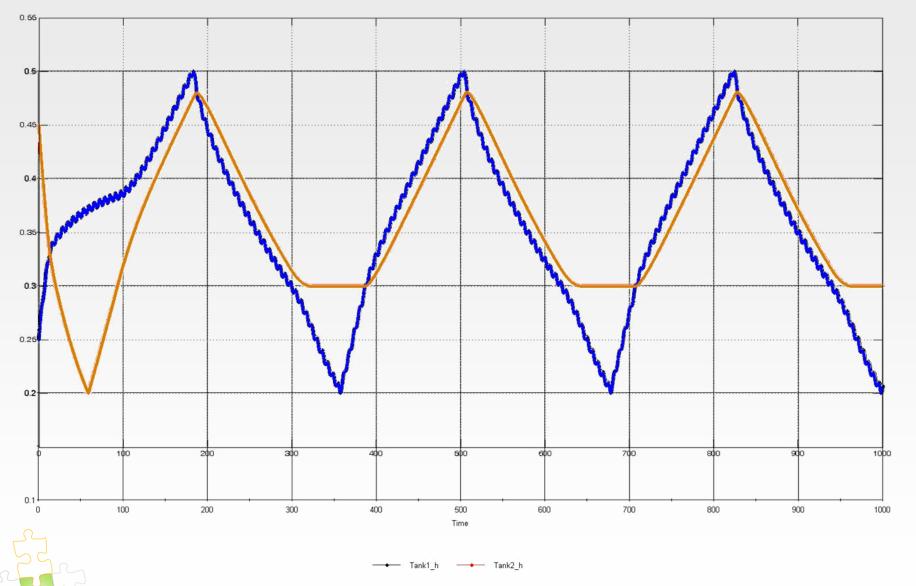
into SFCs

```
model TwoTanks SFC () =
                                                                                                                           Chain
extern var Tanks DOT Tank1 DOT h: cont real
; t_ld
; t u
      mode \ v \ tr1 = when \ l \ u \ now \ do \ (c \ c,t \ rem) := (0, t \ c - (time \ mod \ t \ c)) \ goto \ v \ tr2
     , v_{tr2} = when c_c >= t_{rem} now do (opt_V1,opt_V2,R_V1,R_V2,l_R_V2,l_u,c_c) := (0,0,false,false,true,false,0) goto v_tr3
; Tal
           I( //action "SL V1 End1"
     , V_
           mode v a0 = when R V1 = false and s End1 and Lact V1 End1 SL now do
inter
     (tru
//tra
               // Main structure automaton a str.m
; s_
     (no
               meda
          , V_
               V
      (Tal
; g_
           w
               W
     :: V
; g_
                   mode
                   vi 2= when I par2 and I str2 and not(I strm) now do (s S2,I str2,not finished2):=(true,false,true) goto v s S2
; g_
                   when not(I par2) and I str2 and not(I strm) now do (I str2,not finished2):=(false,false) goto vi 2
; g_
          )|
                   , v s S2= when (run) and I str2 now do (s Start2, s S2,I str2,not finished2):=(true,false,false,true) goto v s Start2
; g_
      mq
                   when not(run) and I str2 now do (I str2,not finished2):=(false,false) goto v s S2
      (Ta
; g_
                   , v_s_Start2= when (Tanks_DOT_Tank2_DOT_h<=t_lower) and I_str2 now do (s_Hei2,s_Start2,l_str2,not_finished2):=(true,false,false,true)
           //a
; g_(
                   goto v s Hei2
           md
               W
     (Ta
                   when (Tanks DOT Tank2 DOT h>t upper) and not(Tanks DOT Tank2 DOT h<=t lower) and I str2 now do
; g_
           wl
      nd
                   (s Low2,s Start2,l str2,not finished2):=(true,false,false,true) goto v s Low2
; g_
           ::V
      (tru
                   when not(Tanks DOT Tank2 DOT h<=t lower or Tanks DOT Tank2 DOT h>t upper) and I str2 now do
; g_
               )|
      OT
                   (not finished2,I str2):=(false,false) goto v s Start2
; g_
      <t
                   , v s Hei2= when (Tanks DOT Tank2 DOT h>t lower) and I str2 now do (s End2,s Hei2,I str2,not finished2):=(true,false,false,true)
; g_
; not
                   goto v s End2
           //a
; SL
                   when not(Tanks DOT Tank2 DOT h>t lower) and I str2 now do (not finished2,I str2):=(false,false) goto v s Hei2
               m
                   , v s Low2= when (Tanks DOT Tank2 DOT h<=t upper) and I str2 now do (s End2,s Low2,I str2,not finished2):=(true,false,false,true)
; t d
               vil
           w
                   goto v s End2
          , V_
                W
//trail
                   when not(Tanks DOT Tank2 DOT h<=t upper) and I str2 now do (not finished2,I str2):=(false,false) goto v s Low2
           W
                   , v s End2 = when (not(run)) and I str2 now do (s FEnd 1,s End2,I par2,I str2,not finished2):=(true,false,false,false,false) goto vi 2 //
; opt
                W
inter<del>n cio</del>
                   last step
intern clo
                   when (run) and not(not(run)) and I str2 now do (s Start2,s End2,I str2,not finished2):=(true,false,false,true) goto v s Start2
                   when not(not(run) or run) and I str2 now do (not finished2,I str2):=(false,false) goto v s End2
                   // last step, return to start (step variable is deactivated by main automaton)
                   :: vi 2
                  )|
               (s
                  //SFC end
               )| // structure automaton
```

Two-tank Example: Simulation Results

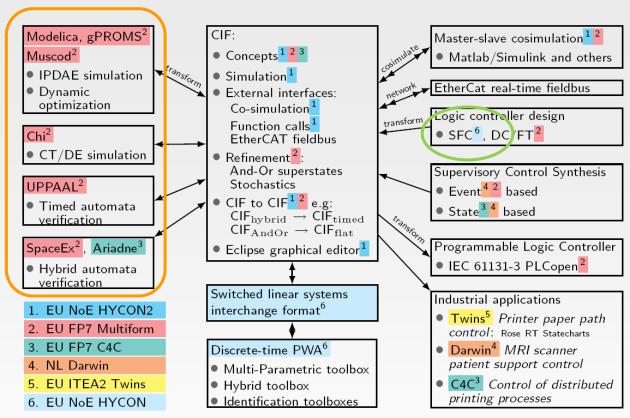


Output Identical





Compositional Interchange Format (CIF)



http://se.wtb.tue.nl/sewiki/cif/start

References:

Fischer, S.; Hüfner, M.; Sonntag, C.; Engell, S.: Systematic Generation of Logic Controllers in a Model-based Multi-formalism Design Environment. Proc. 18th IFAC World Congress, 28.08.-02.09.2011, 12490-12495.

Hendriks, D.; Schiffelers, R.; Hüfner, M.; Sonntag, C.: A Transformation Framework for the Compositional Interchange Format for Hybrid Systems. Proc. 18th IFAC World Congress, 28.08.-02.09.2011, 12509-12514.

Sonntag, C.; Hüfner, M.: On the Connection of Equation- and Automata-based Languages: Transforming the Compositional Interchange Format to Modelica. Proc. 18th IFAC World Congress, 28.08.-02.09.2011, 12515-12520.



Equation-based vs. Automaton-based Formalisms

- Simulation/Solver/Tool options encoded in model code (e.g. EcosimPro, gPROMS)
 - Tool specific options cannot be transformed
 - → Other tools might not find a solution for a difficult initialization problem
- Formal semantics not available -> Transformation not provably correct
- Equation-based models can be more restrictive than automata models
 - E.g. Modelica enforces globally and locally balanced models
 - → Automata models need to be preprocessed
 - Either by flattening of the model
 - Or by rebuilding the automata structure in an equation-based formalism



Summary

- There is a need for efficient model-based support of the design of complex automated systems with trans-level propagation and iteration, and re-use of models
- An all-encompassing mega-tool for the design of complex automated systems is not realistic, so several tools and modeling formalisms must be used in the design process.
- Three different routes to tool and model integration and design support were pursued in MULTIFORM:
 - Model exchange and tool chains via the CIF
 - Direct coupling of tools for testing of specifications
 - Propagation of parameters via the Design Framework



Lessons and Challenges from MULTIFORM

- The CIF and its tool set are stable and relatively mature
- Available under open source licence
- The effort for developing model transformations is high
- Transformations from the CIF in most cases can only be performed for subsets of the models which can be represented in the formalism.
 - A formal specification of the the supported CIF subset of a tool is needed
- It should be possible to trace elements of a model after the transformation
- Model blow-up is not as bad as could be expected



Lessons and Challenges from MULTIFORM (2)

- The CIF is very expressive and well suited for model exchange between automata-based formalisms, but conceptually different from equation-oriented languages (e.g. Modelica, gPROMS, EcosimPro)
 - Possible solution: Use a Modelica subset as an exchange formalism for equation-oriented languages, bridge equation- and automataoriented formalisms via the CIF ↔ Modelica transformation
- Often only some elements of a system are modeled precisely, and these models are formulated in different formalisms (*fragmented* modeling)
 - How can the interdependencies between model fragments be formally described and exploited?

Model ontology needed

- Specification of model formalism expressivity using a common formal vocabulary
- Equipping model artifacts with meta data on their origin(s) → traceability
- Description of relations of partial models to an overall model

